Prolog II

The Notion of Unification

- Unification is when two things "become one"
- Unification is kind of like assignment
- Unification is kind of like equality in algebra
- Unification is mostly like pattern matching
- Example:
 - loves(john, X) can unify with loves(john, mary)
 - and in the process, X gets unified with mary

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Unification I

- Any value can be unified with itself.
 - weather(sunny) = weather(sunny)
- A variable can be unified with another variable.
 - -X = Y
- A variable can be unified with ("instantiated to") any Prolog value.
 - Topic = weather(sunny)

Unification II

- Two different structures can be unified if their constituents can be unified.
 - mother(mary, X) = mother(Y, father(Z))
- A variable can be unified with a structure containing that same variable. This is usually a Bad Idea.
 - -X = f(X)

Unification III

- The explicit unification operator is =
- Unification is symmetric: Steve = father(isaac) means the same as father(isaac) = Steve
- Most unification happens implicitly, as a result of parameter transmission.

Scope of Names

- The scope of a variable is the single clause in which it appears.
- The scope of the "anonymous" ("don't care") variable, __, is itself.
 - loves(_, _) = loves(john, mary)
- A variable that only occurs once in a clause is a useless singleton; you should replace it with the anonymous variable

Writing Prolog Programs

- Suppose the database contains loves(chuck, X) :- female(X), rich(X). female(jane). and we ask who Chuck loves,
- female(X) finds a value for X, say, jane
- rich(X) then tests whether Jane is rich

?- loves(chuck, Woman).

abs(X, X) :- X >= 0.

• A predicate consists of multiple clauses, each of which represents a "case"

Clauses as Cases

$$\begin{split} & \text{grandson}(X,Y) := \text{son}(X,Z), \, \text{son}(Z,Y). \\ & \text{grandson}(X,Y) := \text{son}(X,Z), \, \text{daughter}(Z,Y). \\ & \text{abs}(X,Y) := X < 0, \, Y \text{ is -X}. \end{split}$$

Ordering

- Clauses are always tried in order
- buy(X) :- good(X). buy(X) :- cheap(X).

cheap('Java 2 Complete'). good('Thinking in Java').

• What will buy(X) choose first?

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Ordering II

• Try to handle more specific cases (those having more variables instantiated) first.

```
dislikes(john, bill).
dislikes(john, X) :- rich(X).
dislikes(X, Y) :- loves(X, Z), loves(Z, Y).
```

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Ordering III

- Some "actions" cannot be undone by backtracking over them:
 - write, nl, assert, retract, consult
- Do tests before you do undoable actions:

```
- take(A) :-
    at(A, in_hand),
    write('You\'re already holding it!'),
    nl.
```

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Recursion

- Handle the base cases first
 ancestor(X, Y) :- child(Y, X).
 (X is an ancestor of Y if Y is a child of X.)
- Recur only with a simpler case

```
ancestor(X, Y):-
child(Z, X), ancestor(Z, Y).
(X is an ancestor of Y if Z is a child of X and Z is an ancestor of Y).
```

Case Level

• You can often choose the "level" at which you want cases to be defined.

```
son(isaac, steven).
child(X, Y) :- son(X, Y).
male(isaac).
child(isaac, steven).
son(X, Y) :- male(X), child(X, Y).
```

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Recursive Loops

- Prolog proofs must be tree structured, that is, they may not contain recursive loops.
 - child(X,Y) :- son(X,Y).
 - son(X,Y) :- child(X,Y), male(X).
 - ?- son(isaac, steven). ??? May loop!
- Why? Neither child/2 nor son/2 is atomic

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Cut and Cut-fail

- The cut, !, is a commit point. It commits to:
 - the clause in which it occurs (can't try another)
 - everything up to that point in the clause
- Example:
 - loves(chuck, X) :- female(X), !, rich(X).
 - Chuck loves the *first* female in the database, but only if she is rich.
- Cut-fail, (!, fail), means give up *now* and don't even try for another solution.

What you can't do

- There are no functions, only predicates
- Prolog is programming in logic, therefore there are few control structures
- There are no assignment statements; the *state* of the program is what's in the database

Workarounds I

- There are few control structures in Prolog, BUT...
- You don't need IF because you can use multiple clauses with "tests" in them
- You seldom need loops because you have recursion
- You can, if necessary, construct a "fail loop"

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Workarounds II

- There are no functions, only predicates, BUT...
- A call to a predicate can instantiate variables: female(X) can either
 - look for a value for X that satisfies female(X), or
 - if X already has a value, test whether female(X) can be proved true
- By convention, output variables are put last

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Fail Loops

```
notice_objects_at(Place) :-
    at(X, Place),
    write('There is a '), write(X),
    write(' here.'), nl,
    fail.
notice_objects_at(_).
```

• Use fail loops sparingly, if at all.

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Workarounds III

- There are no assignment statements, BUT...
- the Prolog database keeps track of program state
 - assert(at(fly, bedroom))
 bump_count : retract(count(X)),

Y is X + 1,

assert(count(Y)).

• Don't get carried away and misuse this!

The End