

# Adversarial Search (aka Games)

Chapter 5

# Why study games?

- Interesting, hard problems requiring minimal "initial structure"
- Clear criteria for success
- Study problems involving {hostile, adversarial, competing} agents and uncertainty of interacting with the natural world
- People have used them to assess their intelligence
- Fun, good, easy to understand, PR potential
- Games often define very large search spaces, e.g. chess 35<sup>100</sup> nodes in search tree, 10<sup>40</sup> legal states

# Chess early days



- 1948: Norbert Wiener <u>describes</u> how chess program can work using minimax search with an evaluation function
- 1950: Claude Shannon publishes <u>Programming a</u> <u>Computer for Playing Chess</u>
- 1951: Alan Turing develops *on paper* 1st program capable of playing full chess games (<u>Turochamp</u>)
- 1958: 1st program plays full game on IBM 704 (loses)
- 1962: Kotok & McCarthy (MIT) 1st program to play credibly
- 1967: Greenblatt's Mac Hack Six (MIT) defeats a person in regular tournament play

#### State of the art

- 1979 Backgammon: BKG (CMU) tops world champ
- 1994 Checkers: Chinook is the world champion
- 1997 Chess: IBM Deep Blue beat Gary Kasparov
- 2007 Checkers: solved (it's a draw)
- 2016 Go: AlphaGo beat champion Lee Sedol
- 2017 Poker: CMU's <u>Libratus</u> won \$1.5M from 4 top poker players in 3-week challenge in casino
- 20?? Bridge: Expert <u>bridge programs</u> exist, but no world champions yet

# How can we do it?

# Classical vs. Statistical approach

- We'll look first at the classical approach used from the 1940s to 2010
- Then at newer statistical approached of which AlphaGo is an example
- These share some techniques

# Typical simple case for a game

- 2-person game
- Players alternate moves
- Zero-sum: one player's loss is the other's gain
- **Perfect information**: both players have access to complete information about state of game. No information hidden from either player.
- No chance (e.g., using dice) involved
- Examples: Tic-Tac-Toe, Checkers, Chess, Go, Nim, Othello
- But not: Bridge, Solitaire, Backgammon, Poker, Rock-Paper-Scissors, ...

#### Can we use ...



- Uninformed search?
- Heuristic search?
- Local search?
- Constraint based search?

None of these model the fact that we have an **adversary** ...

# How to play a game

- A way to play such a game is to:
  - -Consider all the legal moves you can make
  - -Compute new position resulting from each move
  - -Evaluate each to determine which is best
  - –Make that move
  - -Wait for your opponent to move and repeat
- Key problems are:
  - -Representing the "board" (i.e., game state)
  - -Generating all legal next boards
  - -Evaluating a position

#### **Evaluation function**

- Evaluation function or static evaluator used to evaluate the "goodness" of a game position

  Contrast with heuristic search, where evaluation function estimates cost from start node to goal passing through given node
- Zero-sum assumption permits single function to describe goodness of board for both players
  - $-\mathbf{f}(\mathbf{n}) >> \mathbf{0}$ : position n good for me; bad for you
  - $-\mathbf{f}(\mathbf{n}) \ll \mathbf{0}$ : position n bad for me; good for you
  - $-\mathbf{f}(\mathbf{n})$  near 0: position n is a neutral position
  - $-\mathbf{f}(\mathbf{n}) = +\mathbf{infinity}$ : win for me
  - $-\mathbf{f}(\mathbf{n}) = -\mathbf{infinity}$ : win for you

# Evaluation function examples

#### • For Tic-Tac-Toe

f(n) = [# my open 3lengths] - [# your open 3lengths] Where 3length is complete row, column or diagonal that has no opponent marks

#### Alan Turing's function for chess

- $-\mathbf{f}(\mathbf{n}) = \mathbf{w}(\mathbf{n})/\mathbf{b}(\mathbf{n})$  where  $\mathbf{w}(\mathbf{n}) = \mathbf{sum}$  of point value of white's pieces and  $\mathbf{b}(\mathbf{n}) = \mathbf{sum}$  of black's
- -Traditional piece values: pawn:1; knight:3; bishop:3; rook:5; queen:9

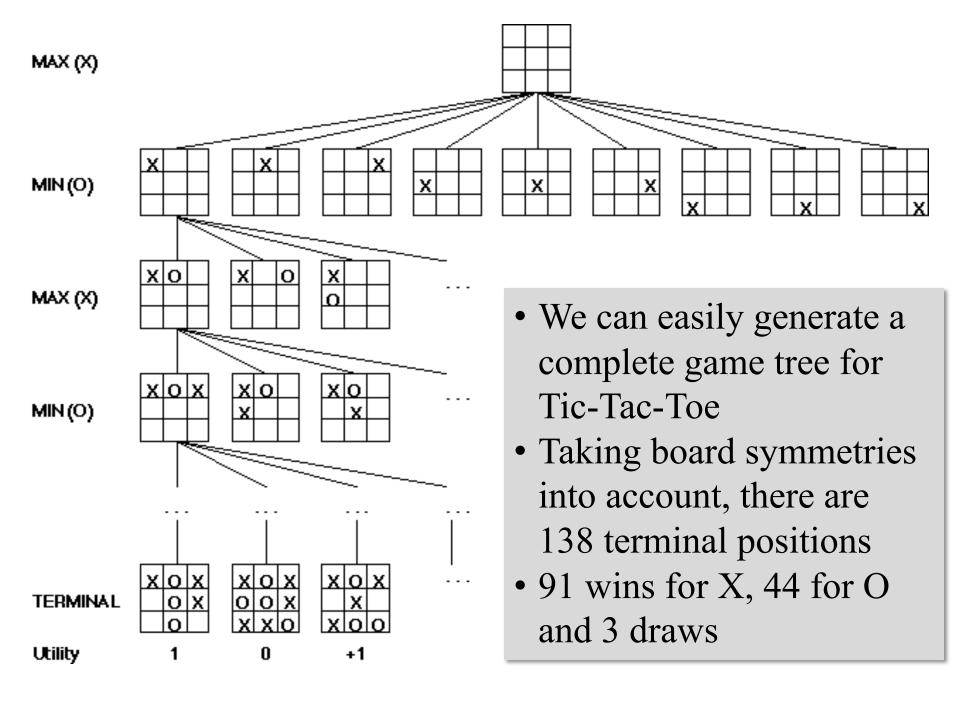
# Evaluation function examples

Most evaluation functions specified as a weighted sum of positive features
 f(n) = w<sub>1</sub>\*feat<sub>1</sub>(n) + w<sub>2</sub>\*feat<sub>2</sub>(n) + ... + w<sub>n</sub>\*feat<sub>k</sub>(n)

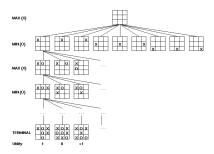
- Example chess features are piece count, piece values, piece placement, squares controlled, etc.
- IBM's chess program <u>Deep Blue</u> (circa 1996) had >8K features in its evaluation function

# But, that's not how people play

- People also use *look ahead* i.e., enumerate actions, consider opponent's possible responses, REPEAT
- Producing a *complete* game tree is only possible for simple games
- So, generate a partial game tree for some number of <u>plys</u>
  - -Move = each player takes a turn
  - -Ply = one player's turn
- What do we do with the game tree?

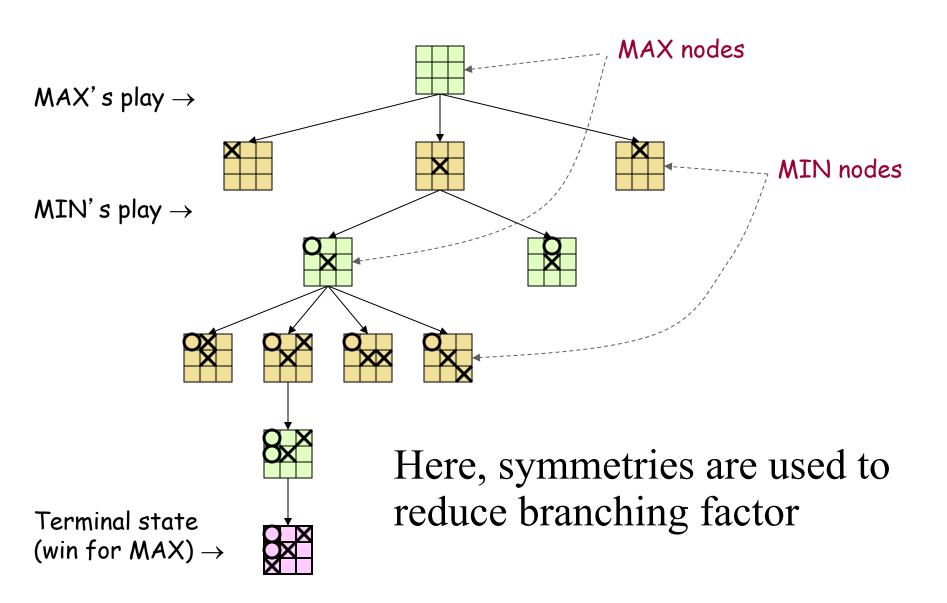


#### Game trees



- Problem spaces for typical games are trees
- Root node is current board configuration; player must decide best single move to make next
- Static evaluator function rates board position f(board):real, >0 for me; <0 for opponent
- Arcs represent possible legal moves for a player
- If my turn to move, then root is labeled a "MAX" node; otherwise it's a "MIN" node
- Each tree level's nodes are all MAX or all MIN; nodes at level i are of opposite kind from those at level i+1

#### Game Tree for Tic-Tac-Toe



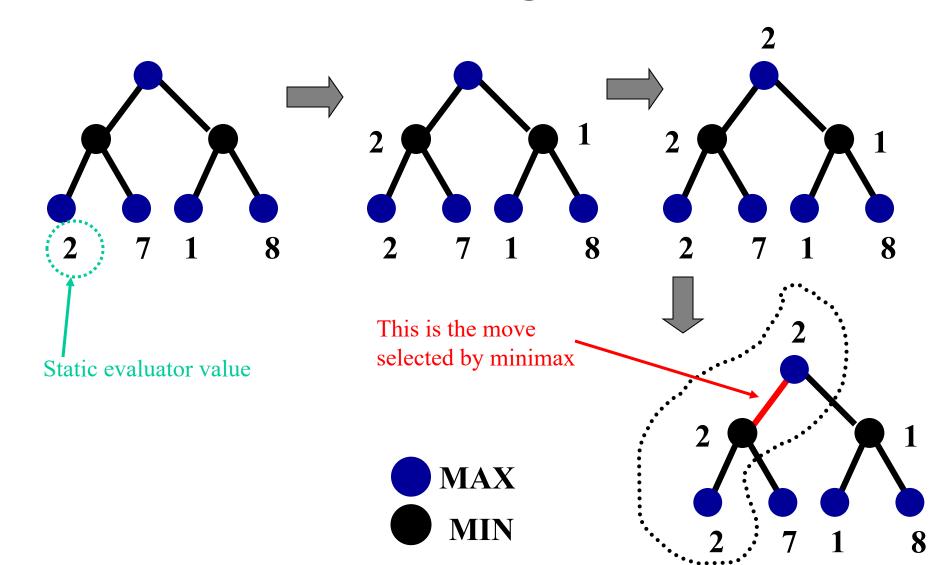
# Minimax procedure

- Create MAX node with current board configuration
- Expand nodes to some **depth** (a.k.a. **plys**) of lookahead in game
- Apply evaluation function at each leaf node
- *Back up* values for each non-leaf node until value is computed for the root node
  - At MIN nodes: value is **minimum** of children's values
  - At MAX nodes: value is **maximum** of children's values
- Choose move to child node whose backed-up value determined value at root

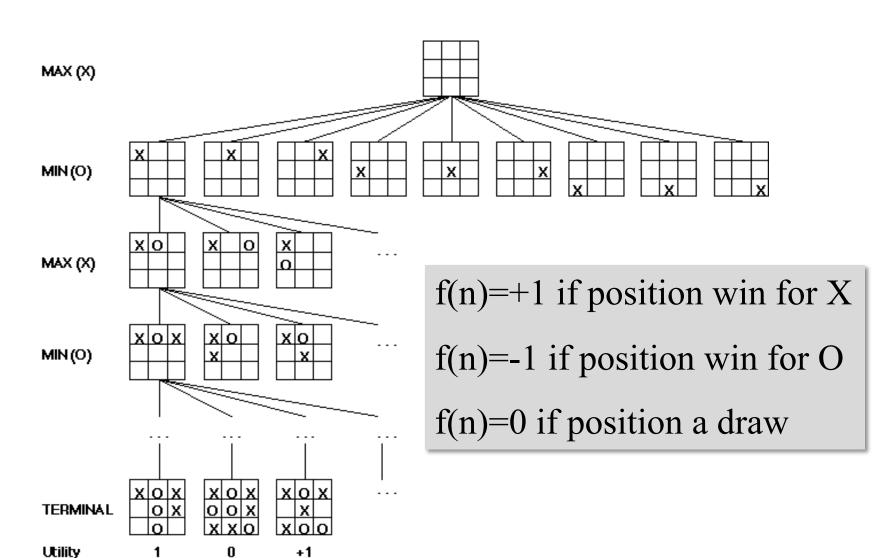
#### Minimax theorem

- Intuition: assume your opponent is at least as smart as you and play accordingly
  - −If she's not, you can only do better!
- Von Neumann, J: Zur Theorie der Gesellschaftsspiele Math. Annalen. **100** (1928) 295-320
  - For every 2-person, 0-sum game with finite strategies, there is a value V and a mixed strategy for each player, such that (a) given player 2's strategy, best payoff possible for player 1 is V, and (b) given player 1's strategy, best payoff possible for player 2 is –V.
- You can think of this as:
  - -Minimizing your maximum possible loss
  - -Maximizing your minimum possible gain

# **Minimax Algorithm**



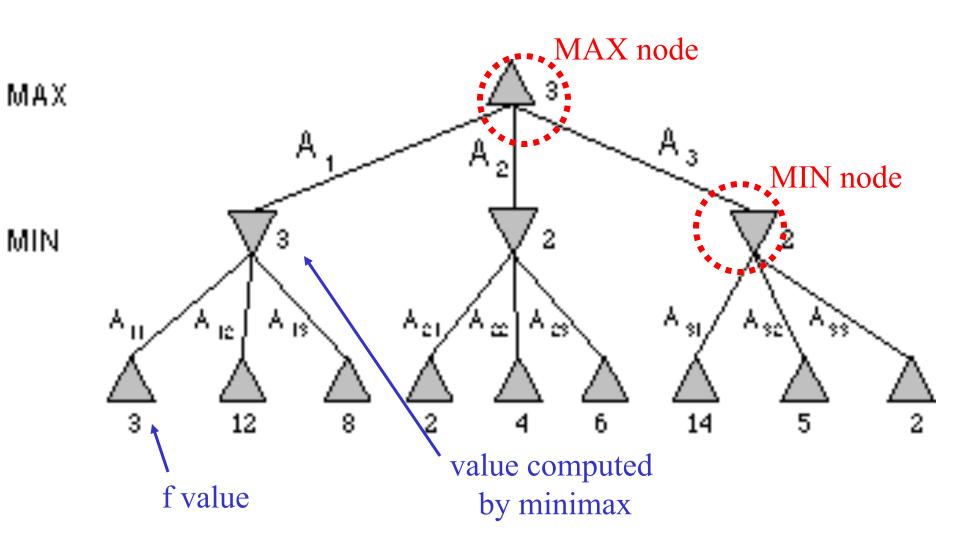
#### Partial Game Tree for Tic-Tac-Toe



#### Why backed-up values?

- Why not just use a good static evaluator metric on immediate children
- Intuition: if metric is good, doing look ahead and backing up values with Minimax should be better
- Non-leaf node N's backed-up value is value of best state MAX can reach at depth **h** if MIN plays *well* 
  - "plays well": same criterion as MAX applies to itself
- If e is good, then backed-up value is better estimate of STATE(N) goodness than e(STATE(N))
- Use lookahead horizon h because time to choose move is limited

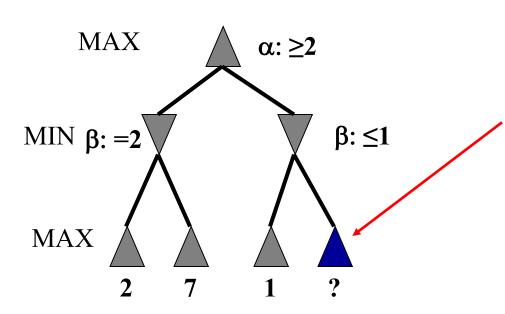
#### **Minimax Tree**



# Is that all there is to simple games?

# Alpha-beta pruning

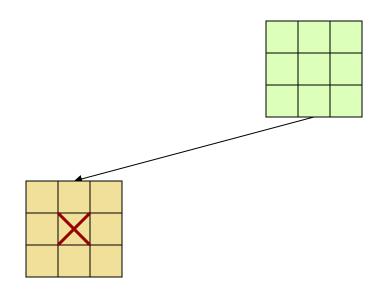
- Improve performance of the minimax algorithm through <u>alpha-beta pruning</u>
- "If you have an idea that is surely bad, don't take the time to see how truly awful it is"-Pat Winston (MIT)

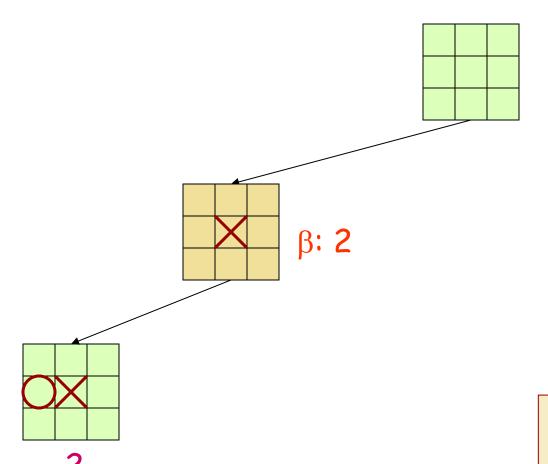


- We don't need to compute the value at this node
- No matter what it is, it can't affect value of the root node

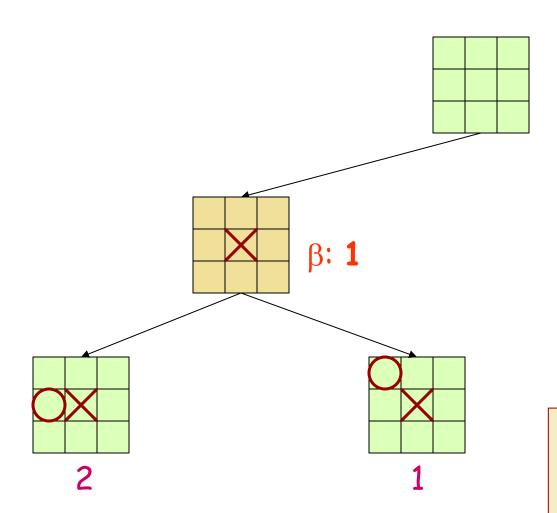
# Alpha-beta pruning

- Traverse search tree in depth-first order
- At MAX node n, alpha(n) = max value found so far Alpha values start at  $-\infty$  and only increase
- At MIN node n, beta(n) = min value found so far Beta values start at  $+\infty$  and only decrease
- **Beta cutoff**: stop search below MAX node N (i.e., don't examine more descendants) if alpha(N) >= beta(i) for some MIN node ancestor i of N
- Alpha cutoff: stop search below MIN node N if beta(N)<=alpha(i) for a MAX node anceastor i of N

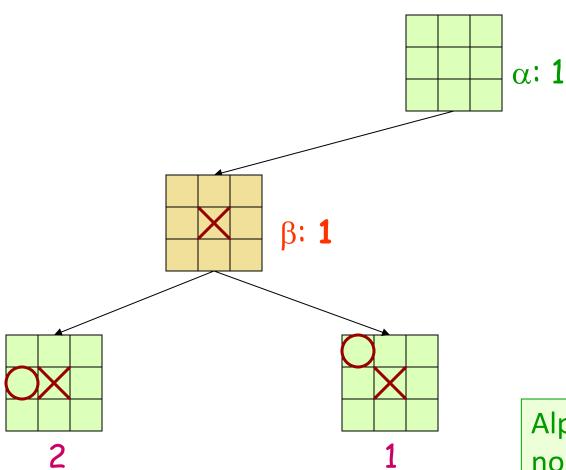




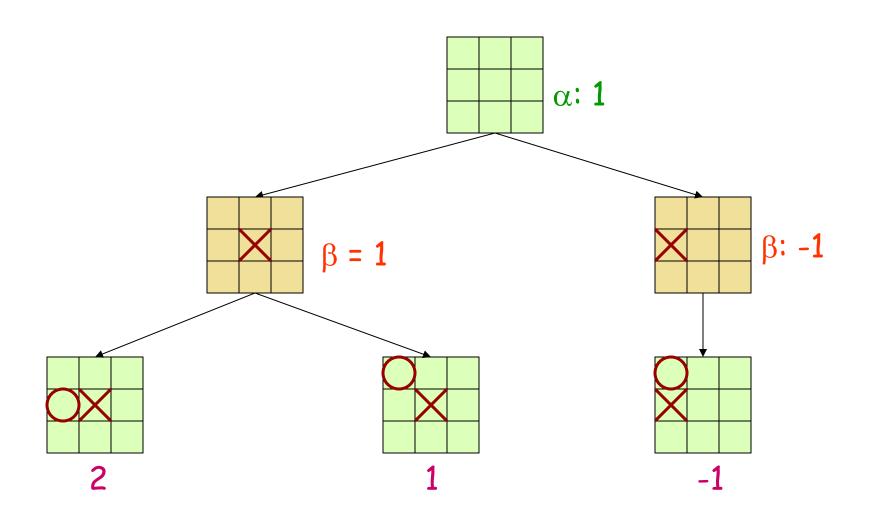
Beta value of a MIN node is **upper** bound on final backed-up value; it can never increase

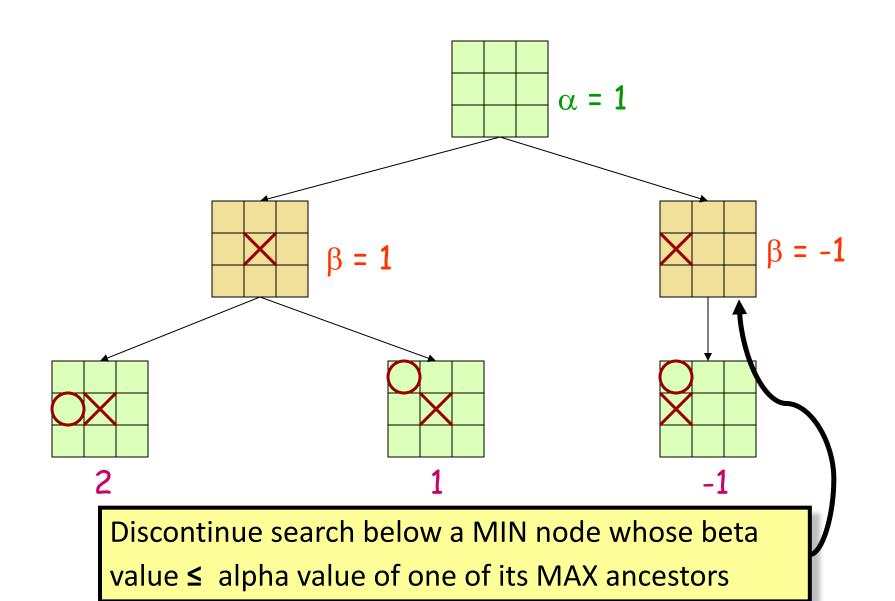


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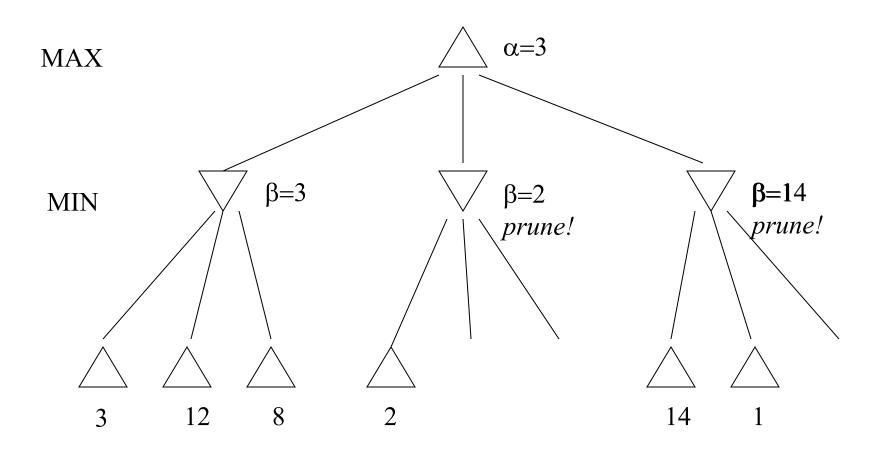


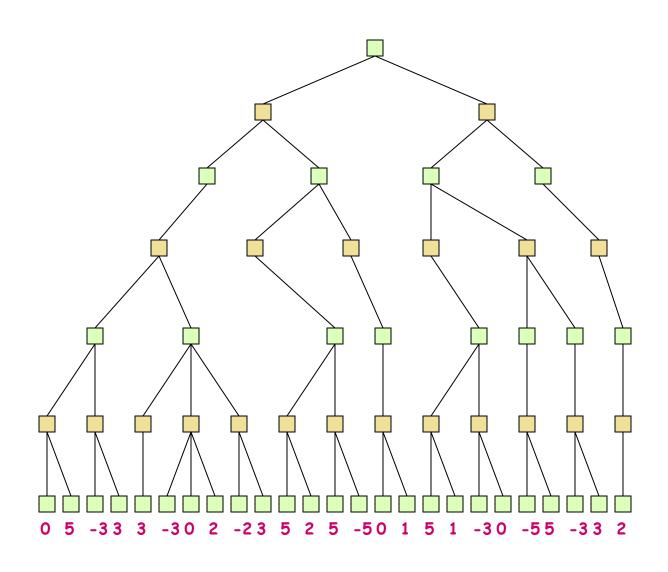
Alpha value of MAX node is **lower** bound on final backed-up value; it can never decrease

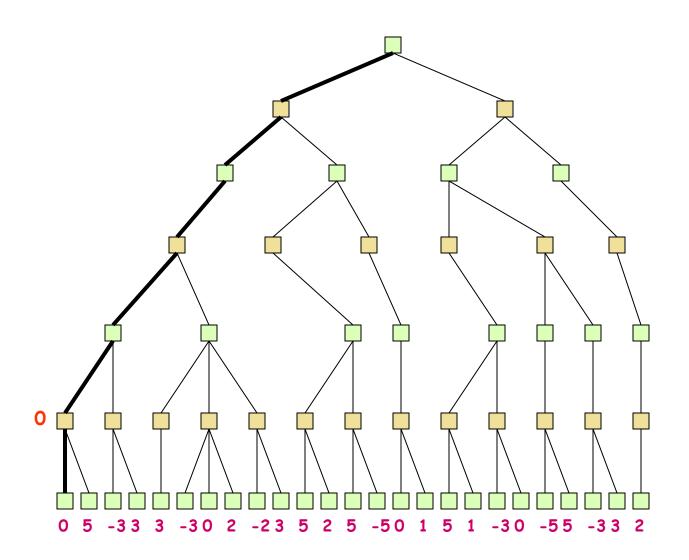


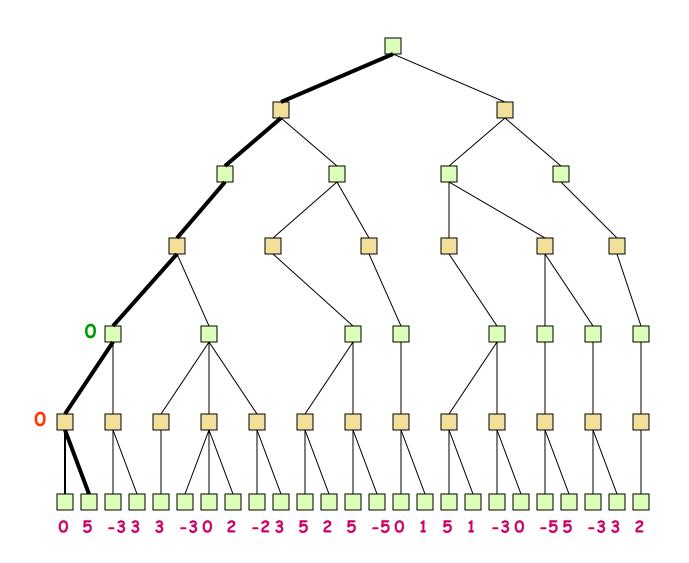


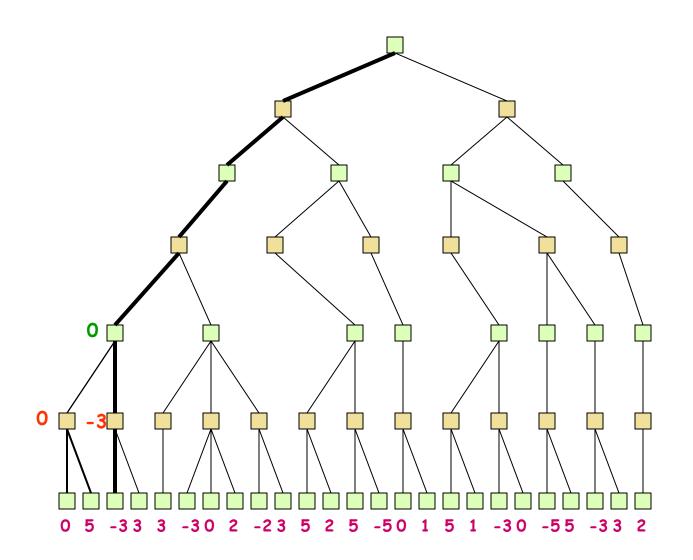
# Another alpha-beta example

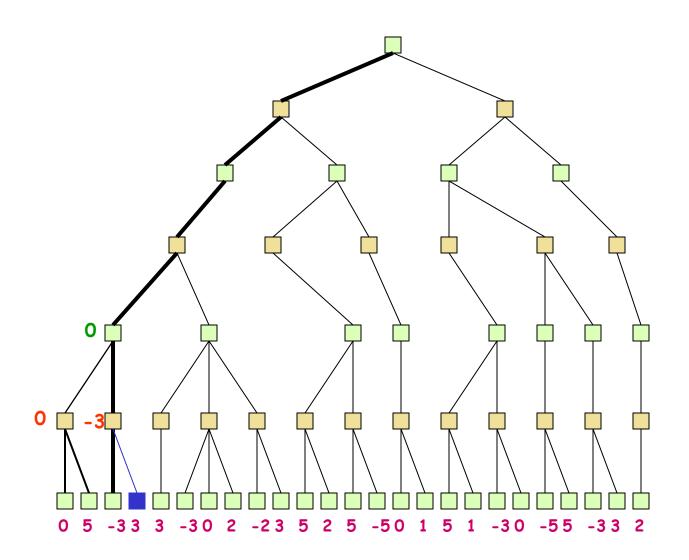


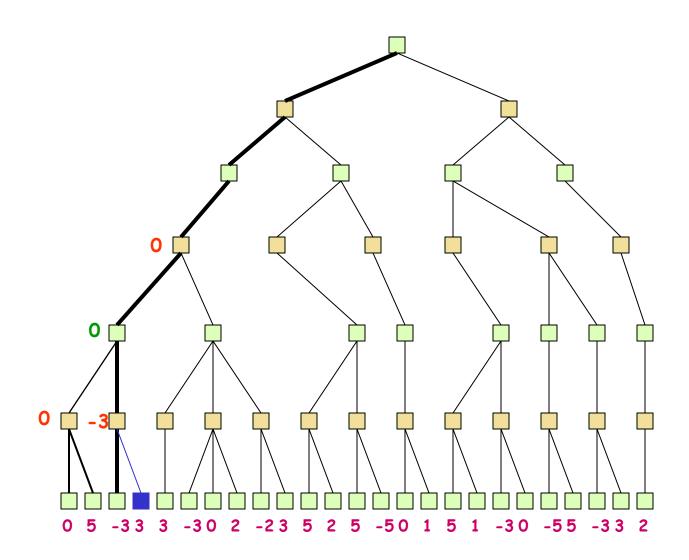


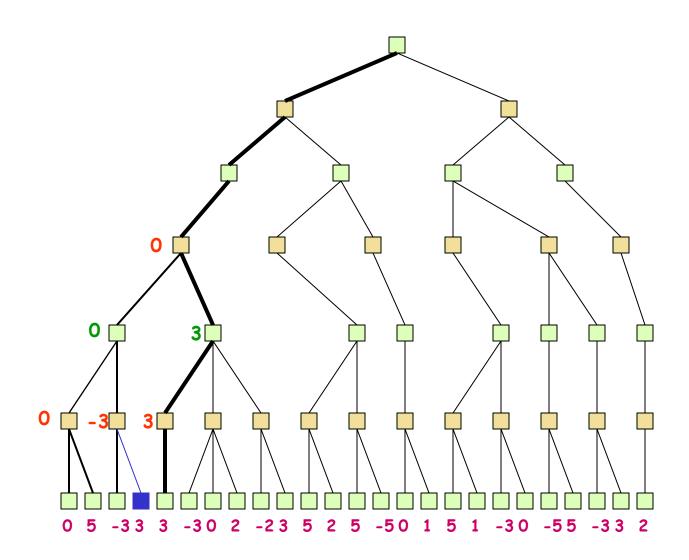


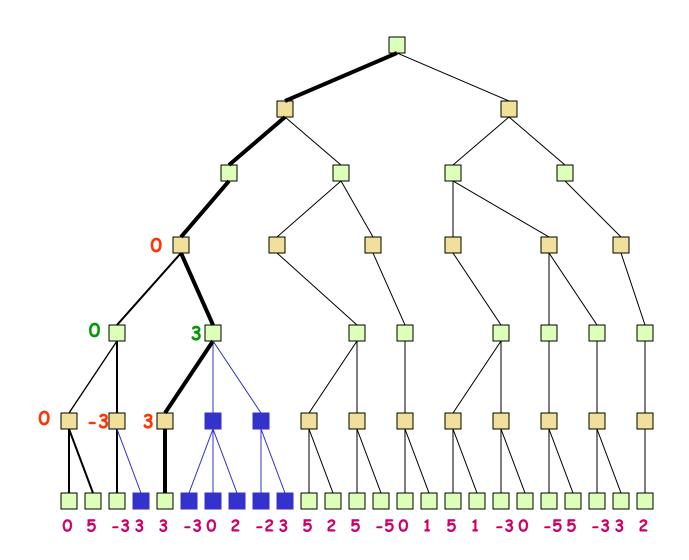


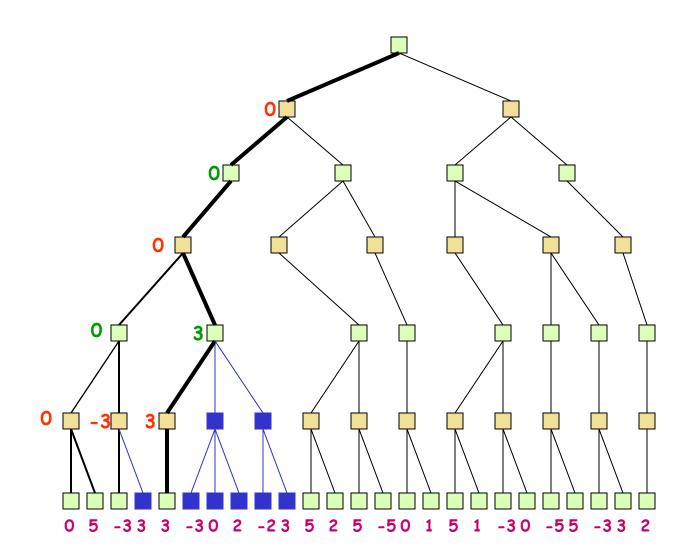


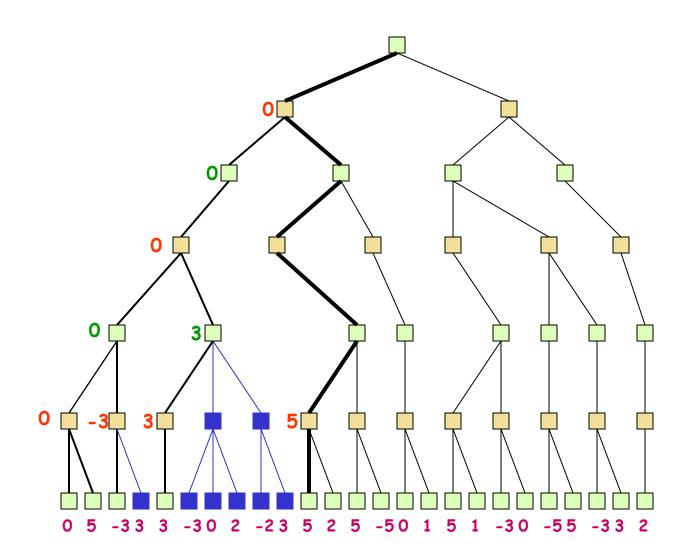


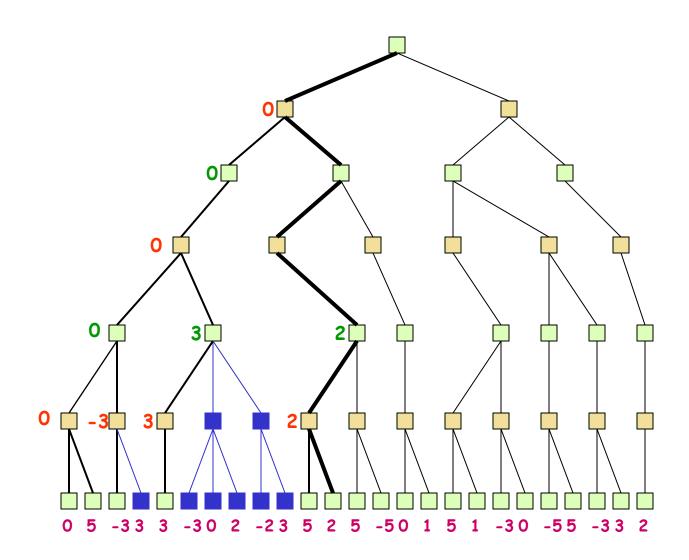


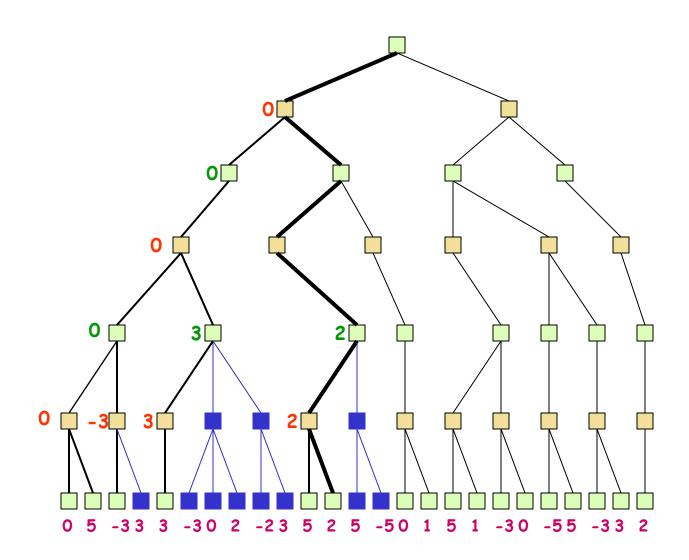


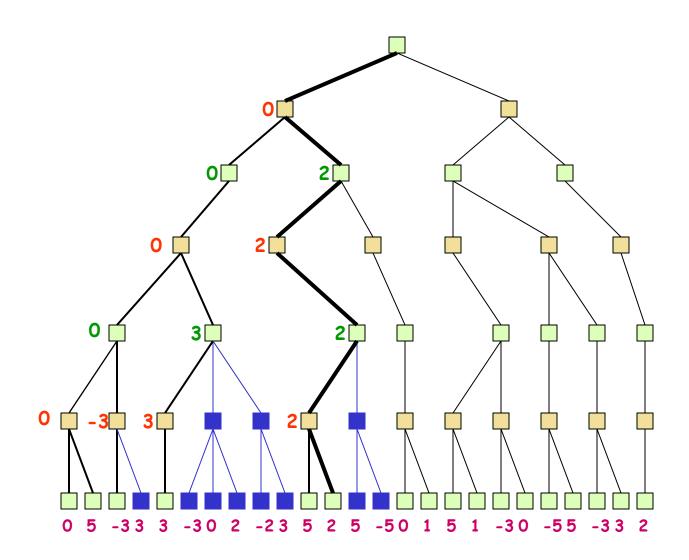


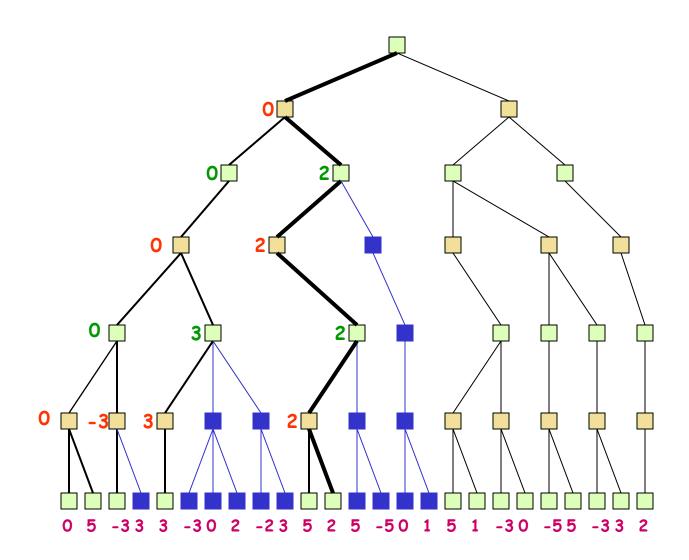


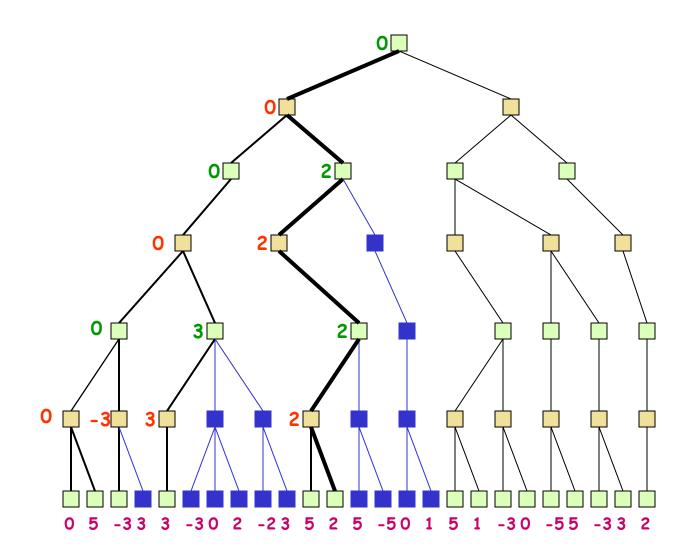


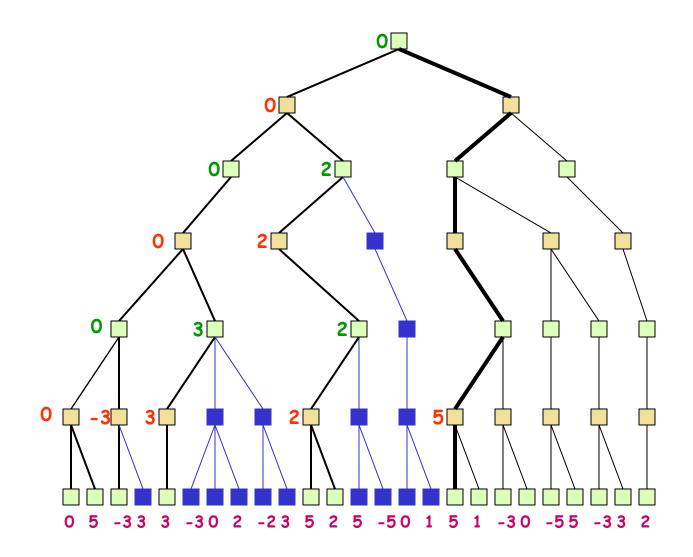


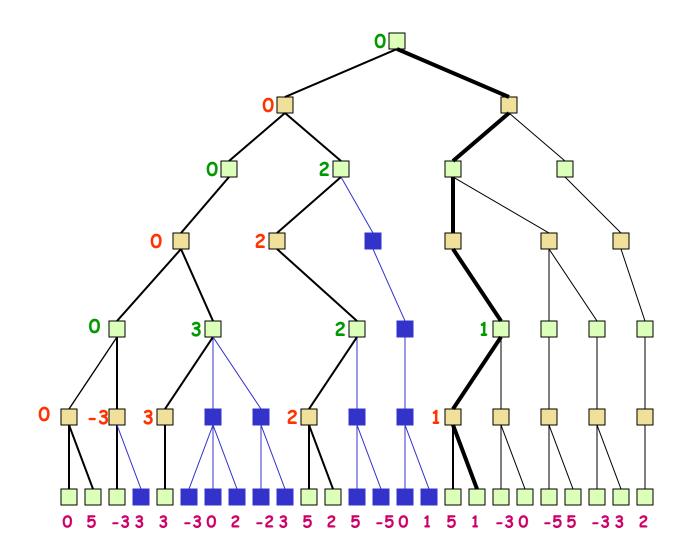


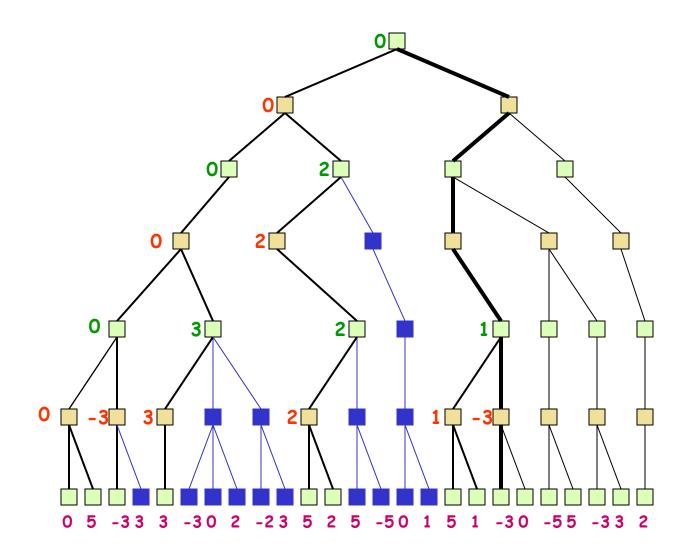


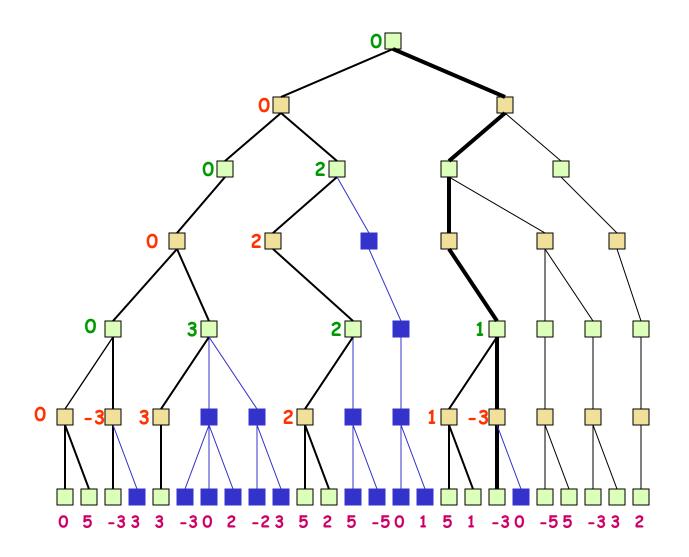


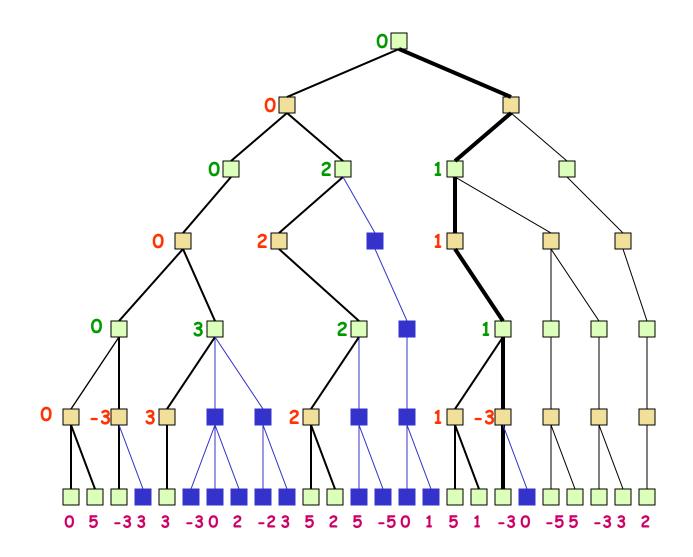


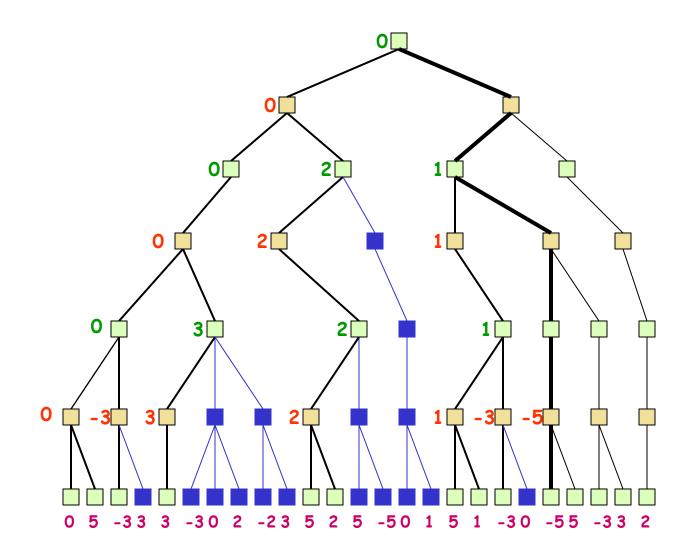


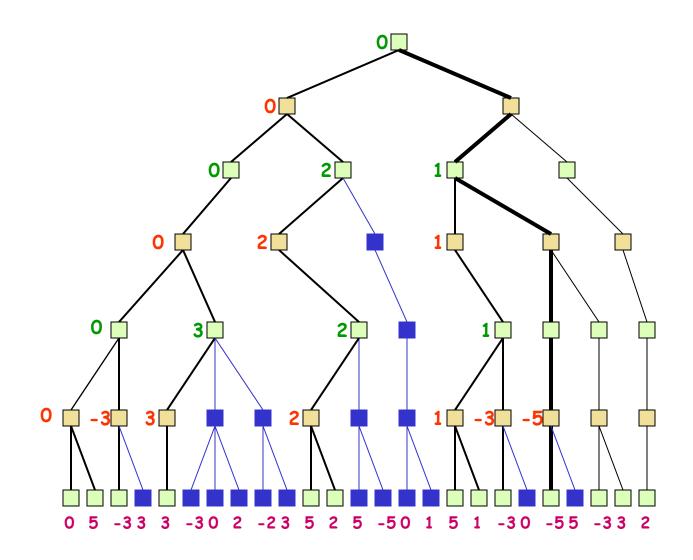


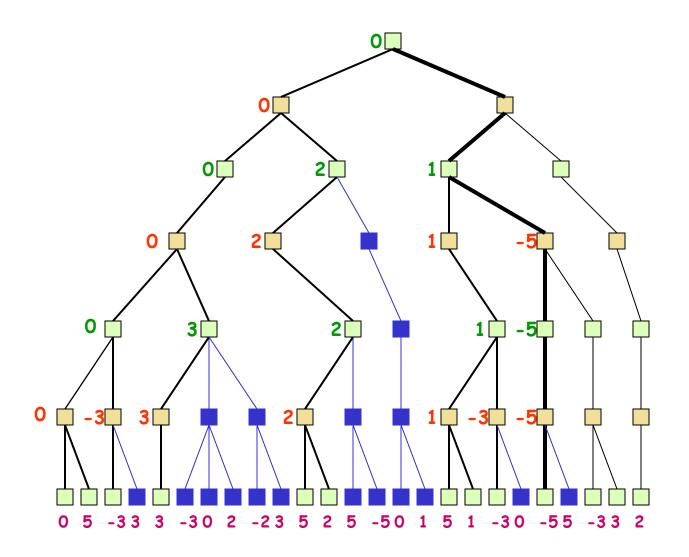


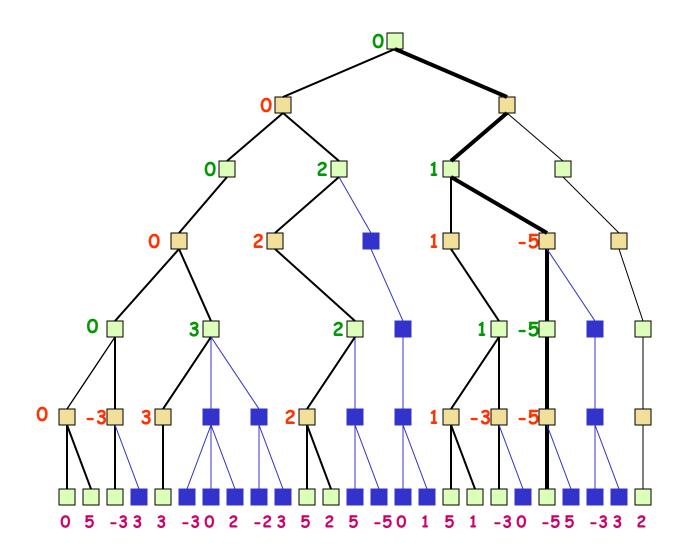


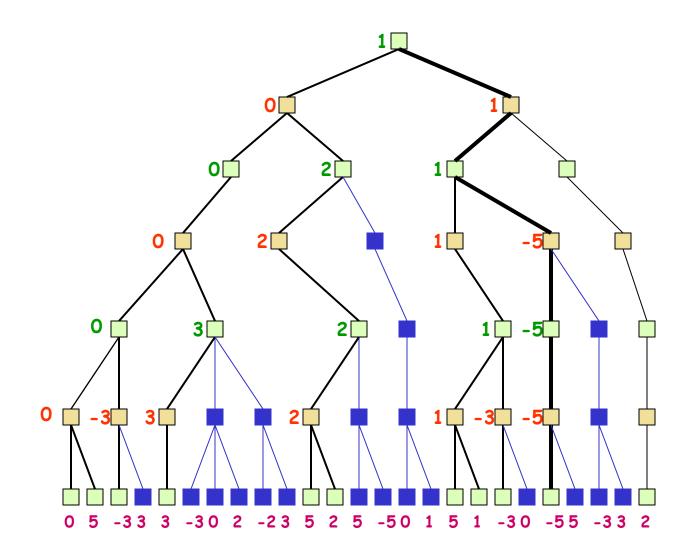


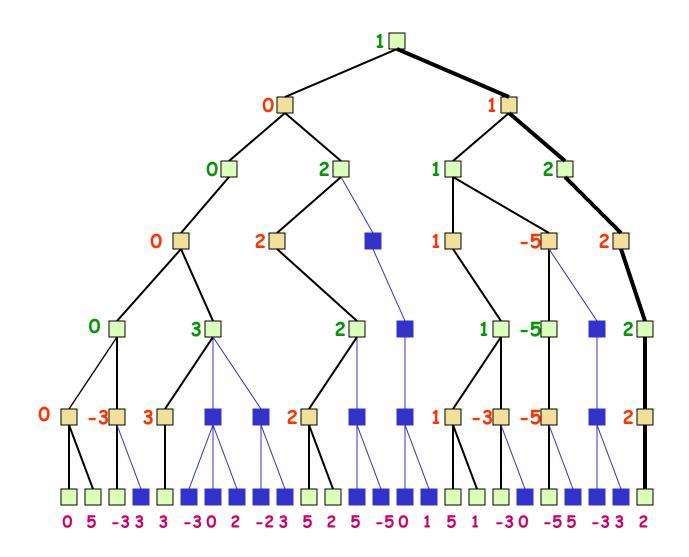




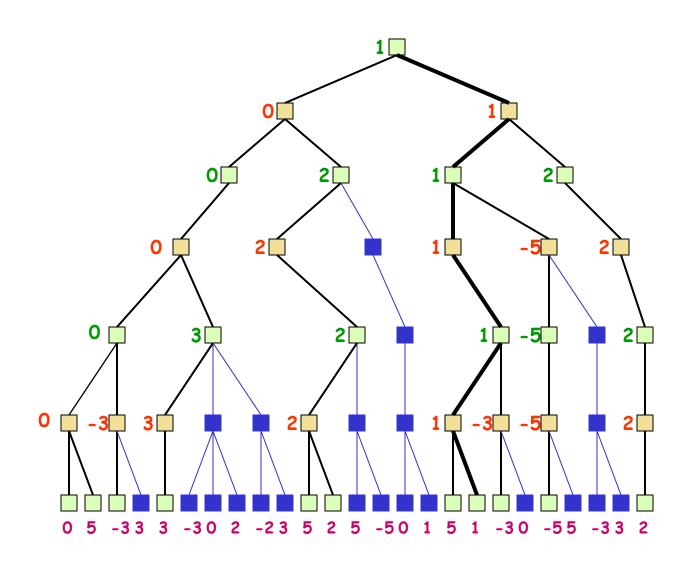








## With alpha-beta we avoided computing a static evaluation metric for 14 of the 25 leaf nodes



## Effectiveness of alpha-beta

- Alpha-beta guaranteed to compute same value for root node as minimax, but with  $\leq$  computation
- Worst case: no pruning, examine b<sup>d</sup> leaf nodes, where nodes have b children & d-ply search is done
- Best case: examine only (2b)<sup>d/2</sup> leaf nodes
  - You can search twice as deep as minimax!
  - -Occurs if each player's best move is 1st alternative
- In <u>Deep Blue</u>, alpha-beta pruning reduced effective branching factor from ~35 to ~6

## Many other improvements

- Adaptive horizon + iterative deepening
- Extended search: retain k>1 best paths (not just one) extend tree at greater depth below their leaf nodes to help dealing with "horizon effect"
- Singular extension: If move is obviously better than others in node at horizon h, expand it
- Use <u>transposition tables</u> to deal with repeated states