UMBC CMSC 671 Midterm Exam
25 October 2015

Write your answers on this exam, which is closed book and consists of seven problems, summing to 100 points. You have the entire class period, seventy-five minutes, to work on this exam. Good luck.

1. True/False [10 points]

Circle either T or an F in the space before each statement to indicate whether the statement is true or false. If you think the answer is simultaneously true and false, quit while you are ahead. There is no penalty for incorrect answers but then, there are no points for incorrect answers either.

T F In a stochastic environment, the next state is completely determined by the agent’s action.  
FALSE. Stochastic environments involve an element of chance.

T F Hill climbing search algorithms only work for search spaces that are two-dimensional or have solution-preserving projections onto two-dimensions.  
FALSE

T F If search heuristic \( H_1(s) \) is admissible and heuristic \( H_2(s) \) is not admissible, then \( H_3(s) = \min(H_1(s), H_2(s)) \) will be admissible.  
TRUE since \( H_3(s) \leq H_1(s) \leq H^*(s) \)

T F In A* search, the first path to the goal which is added to the fringe will always be optimal.  
FALSE. The first path removed from the fringe is optimal with an admissible heuristic.

T F If a Constraint Satisfaction Problem (CSP) is arc consistent, it can be solved without backtracking.  
FALSE. Arc consistency may not determine the entire solution.

T F Deciding if a CSP is consistent is, in general, NP-hard.  
TRUE

T F Arc consistency is a more powerful constraint propagation algorithm than forward checking.  
TRUE

T F Alpha-beta pruning can alter the computed minimax value of the root of a game search tree.  
FALSE. Alpha-beta pruning only speeds up computation; it does not change the answer.

T F Minimax with alpha-beta pruning on a game tree which is traversed from left to right will never prune the leftmost branch.  
TRUE. There are no alternatives to motivate pruning at the left-most branch of a game tree.

T F Chess, checkers and go are examples of games that have a partially observable environment.  
FALSE. The entire state is visible in the board for these games.
2. Search I [20]

Assume the following search graph, where S is the start node and G1 and G2 are goal nodes. Arcs are labeled with the cost of traversing them and the estimated cost to a goal is reported inside nodes.

For each of the search strategies below, show which goal state is reached (if any) and list, in order, the states expanded. (Recall that for algorithm A, a state is expanded when it is removed from the OPEN list and its successors added to the queue). When all else is equal, nodes should be expanded in alphabetical order.

A difficulty with this problem is that the exact version of the algorithms to be used is not specified. I had in mind, and based my answers on, the generic graph-search algorithm discussed in class. However, I suspect that not all of the students did this, so the problem is under-defined. Another issue is that there is a difference between a node being added to the graph (or tree) and it being expanded -- these happen at different times. The question asks about which nodes are expanded.

**Depth first [5]**

<table>
<thead>
<tr>
<th>goal found</th>
<th>G2</th>
</tr>
</thead>
<tbody>
<tr>
<td>states expanded</td>
<td>S B E</td>
</tr>
</tbody>
</table>

- My (preferred) answer is that it reaches G2 after expanding (S,B,E). My assumption is that the graph-based DFS algorithm avoids loops and knows to stop as soon as a goal appears on the fringe.
- A possible answer might be that it loops, which would happen if you assume that (i) the algorithm is not graph-based and (ii) goal states are not recognized until chosen for expansion.
- Another possible answer might be it selects G2 and expands (S,B,E,D,A), which happens if you assume that (i) the algorithm is graph-based so avoids loops and (ii) goal states are not recognized until chosen for expansion.
### Breadth first [5]

<table>
<thead>
<tr>
<th>goal found</th>
<th>G1</th>
</tr>
</thead>
<tbody>
<tr>
<td>states expanded</td>
<td>S B C</td>
</tr>
</tbody>
</table>

My answer of reaching G2 after expanding (S,B,C) is correct assuming goals are discovered as soon as they are added to the search fringe. The nodes E and F are added to the graph, but not expanded.

### Hill Climbing [5] (using the h function only)

<table>
<thead>
<tr>
<th>goal found</th>
<th>G1</th>
</tr>
</thead>
<tbody>
<tr>
<td>states expanded</td>
<td>S B F</td>
</tr>
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</table>

### Algorithm A (Graph Search) [5]

<table>
<thead>
<tr>
<th>goal found</th>
<th>G2</th>
</tr>
</thead>
<tbody>
<tr>
<td>states expanded</td>
<td>S B E D</td>
</tr>
</tbody>
</table>
3. Constraint Satisfaction [20]

You are scheduling five lower-level CMSC courses for the spring that have three professors available to teach them. A professor can teach more than one course, but only of the times don’t overlap. The courses and the times when they meet are:

- Course 1: CMSC201, 8:00-9:00am
- Course 2: CMSC202, 8:30-9:30am
- Course 3: CMSC203, 9:00-10:00am
- Course 4: CMSC331, 9:00-10:00am
- Course 5: CMSC341, 9:30-10:30am

The professors are:

- Prof. A, who is available to teach Courses 3 and 4
- Prof. B, who is available to teach Courses 2, 3, 4 and 5
- Prof. C, who is available to teach Courses 1, 2, 3, 4 and 5

(a) [5] Formulate this as a CSP problem with one variable per course and give the initial domain (i.e., set of possible values) after applying the unary constraints (i.e., which Courses a professor can teach).

<table>
<thead>
<tr>
<th>Variable</th>
<th>Initial domain after applying unary constraints</th>
</tr>
</thead>
<tbody>
<tr>
<td>C1</td>
<td>C</td>
</tr>
<tr>
<td>C2</td>
<td>B, C</td>
</tr>
<tr>
<td>C3</td>
<td>A, B, C</td>
</tr>
<tr>
<td>C4</td>
<td>A, B, C</td>
</tr>
<tr>
<td>C5</td>
<td>B, C</td>
</tr>
</tbody>
</table>

(b) [5] List all of the constraints between the variables

- C1 ≠ C2
- C2 ≠ C3
- C3 ≠ C4
- C4 ≠ C5
- C2 ≠ C4
- C3 ≠ C5

(3) [5] Show the domains of the variables after running the arc-consistency algorithm

<table>
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</tr>
</thead>
<tbody>
<tr>
<td>C1</td>
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</tr>
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<td>C2</td>
<td>B</td>
</tr>
<tr>
<td>C3</td>
<td>A, C</td>
</tr>
<tr>
<td>C4</td>
<td>A, C</td>
</tr>
<tr>
<td>C5</td>
<td>B, C</td>
</tr>
</tbody>
</table>

Note: C5 can’t be C, but arc consistency doesn’t rule it out

(4) [5] Give one solution to this CSP by showing an assignment for each variable

Two possible answers:
- C1 = C; C2 = B; C3 = C; C4 = A; C5 = B
- C1 = C; C2 = B; C3 = A; C4 = C; C5 = B

Consider the following simple game tree in which the root node is the maximizing player. The values that will be returned by the static evaluator are shown for the leaf nodes.

![Game Tree](image)

(a) [5] What are the values that would be backed up by the standard minimax procedure (i.e., operating in a left-to-right traversal of the tree) in the three blank nodes.
- root: [ -4 ]; level2 left: [ -5 ], level2 center: [ -4 ]; level2 right: [ -9 ]

(b) [5] Which move does minimax choose: [ ] leftmost; [ X ] center; [ ] rightmost

(c) [5] Put an X through all leaf nodes whose score need not be computed if the standard alpha-beta version of minimax is used. **Only the right most leaf node (with value 9) need not be evaluated**

5. Propositional logic 1 [5]

A propositional sentence is well formed if it follows the syntax of propositional logic, satisfiable if there is a way to assign true or false to each of its variables that makes the value of the overall sentence true, unsatisfiable if there is no way to assign true or false to its variables that makes the sentence true, and valid if it is always true no matter what values its variables are assigned.

Circle all of the following that are true: The sentence \((P \rightarrow Q) \land Q \rightarrow P\) is

(a) well formed
(b) valid
(c) satisfiable
(d) unsatisfiable

The sentence is (a) well formed; (c) satisfiable. **A model for the sentence is (P=True, Q=False) so it is satisfiable. The model (P=False, Q=True) makes the sentence false, so it is not valid.**
6. Propositional Logic 2 [15]

Let these propositional symbols have the following meaning:

<table>
<thead>
<tr>
<th>symbol</th>
<th>meaning</th>
</tr>
</thead>
<tbody>
<tr>
<td>A</td>
<td>The patient was in an accident</td>
</tr>
<tr>
<td>S</td>
<td>The patient is sick</td>
</tr>
<tr>
<td>I</td>
<td>The patient is injured</td>
</tr>
<tr>
<td>D</td>
<td>The patient needs to see a doctor</td>
</tr>
</tbody>
</table>

Express each of the following English sentences in propositional logic using the symbols \(\land, \lor, \neg, \rightarrow\) and \(\leftrightarrow\) for the logical connectives and, or, not and implies and ‘if and only if’.

(a) [5] The patient was in a car accident, but is not injured

\[ A \land \neg I \]

(b) [5] The patient needs to see a doctor if he is sick or injured

\[ S \lor I \rightarrow D \]

(c) [5] If the patient was not in an accident and is not sick, then he does not need to see a doctor

\[ \neg A \land \neg S \rightarrow \neg D \]

7. English to logic (15)

Translate the following statements into a single sentence in first order logic, choosing appropriate predicates and functions. Avoid redundancy. Use the logical connectives and quantifiers \(\forall\) and \(\exists\) and numeric relations \(<\) and \(>\).

(a) Every boy is a human who is male and whose age is less than 16 and every male human who is less than 16 is a boy

\[ \forall x \text{ boy}(x) \leftrightarrow \text{human}(x) \land \text{male}(x) \land \text{age}(x) < 16 \]

(b) Some people like every vegetable

\[ \exists x \forall y \text{ person}(x) \land \text{vegetable}(y) \rightarrow \text{likes}(x,y) \]

(c) There is no vegetable that is liked by every person

Two simple sentences, either of which is acceptable:

\[ \neg \exists x \forall y \text{ vegetable}(x) \land \text{person}(y) \rightarrow \text{likes}(x,y) \]
\[ \forall x \exists y \text{ vegetable}(x) \land \text{person}(y) \rightarrow \neg \text{likes}(x,y) \]