# Preventing Stack-based Buffer Overflows

CMSC 426 - Computer Security

## Outline

- Buffer Overflows Today
- Safe Programming
- Protection Technologies

## **Buffer Overflows Today**

- Vulnerabilities persist despite protection
- A <u>vulnerability announcement</u> from 5 February 2018
- You can find plenty of published buffer overflow vulnerabilities:
  - National Vulnerability Database
  - US Computer Emergency Response

# Safe Programming

• Avoid unsafe functions. Use safe alternatives; for example:

Unsafe	Safe	Description
strcpy	strncpy	Copy a string
sprintf	snprintf	Formatted print to a string
strcat	strncat	Concatenate strings
gets	fgets	Read string until newline or EOF

<pre>#include <stdio.h></stdio.h></pre>	Will use gets()
#define BUFLEN 20	james:it430 cmarron\$ gcc gets_fgets
	james:it430 cmarron\$ ./a.out
main(int argc, char *argv[]) {	warning: this program uses gets(), which is unsafe.
	***************************************
char name[BUFLEN];	***************************************
	***************************************
if (argc == 1) {	******
	***************************************
gets(name);	***************************************
	***************************************
} else {	xxxxxxxx Will use fgets()
	Segmentation fault: 11
fgets(name, BUFLEN, stdin);	james:it430 cmarron\$ ./a.out x
	***************************************
}	***************************************
	***************************************
<pre>printf("%s\n", name);</pre>	XXXXXXXX
	***************************************
}	james:it430 cmarron\$

#### More Safe Programming

- Always check user input ... but it's hard to catch every possible user error.
- There are tools that help identify unsafe programming, memory leaks, etc., e.g. *Rational Purify*, *Insure++*, *Valgrind* (Open Source).

## Random Canaries

- Run-time protection, but may be provided by the compiler (e.g. GCC's stack protector).
  - Random value (the *canary*) on stack between local variables and return address
  - Change in canary value indicates a buffer overflow stop execution
  - Inflicts some additional overhead on code

## Stack with Canary

Return Address	0xbffff498
Random Canary	0xa18a6f6c
	0x00000000
	0x0000000
<b>Start of</b> char check[64]	0x00000000

## Stack Code Execution

- A standard buffer overflow attack involves executing shellcode on the stack
- One means of protection is to disallow execution of code on the stack
- By default, Linux **disallows** stack execution, but it can be enabled by flipping a bit in the ELF header:

-z execstack option with GCC or execstack utility

# Address Space Layout Randomization

- Process address space re-arranged randomly
- More difficult for attacker to determine addresses of shellcode, libraries, etc.
  - E.g. return-to-libc requires exact address; ASLR makes this difficult
- Implemented to one degree or another in all major OSs (including mobile OSs).

# Linux ASLR

- For Ubuntu, <u>ASLR</u> has been turned-on by default, at least to some degree, since 6.06.
- Part of the OS; not controlled by the compiler.
- Disabling ASLR requires root access:

sysctl -w kernel.randomize\_va\_space=0

## Linux Summary

- Suppose you wanted to write a buffer overflow attack for Linux (Ubuntu, at least); you would have to disable three protections...
  - Turn-off stack protection with the -fno-stackprotector compile option.
  - 2. Allow for an executable stack with the -z execstack linker option.
  - 3. Turn-off ASLR with the system command sysctl -w kernel.randomize va space=0.

## Return Address Defense

• Encrypt the return addresses; decrypted when the address is loaded into a register

PointGuard system XORs return address with 32-bit "key"; decodes at return time.

• Save copy of return address in protected memory; at return time, compare address on stack with saved address

E.g. Stackshield or Return Address Defender

#### **Next time: Malicious Software**