## Functional Programming in Scheme and Lisp



## Overview

- In a functional programming language, functions are first class objects
- You can create them, put them in data structures, compose them, specialize them, apply them to arguments, etc.
- We'll look at how functional programming things are done in Lisp


## eval

- Remember: Lisp code is just an s-expression
- You can call Lisp's evaluation process with the eval function
$>$ (define s (list 'cadr ' ' (one two three)))
>s
(cadr ' (one two three))
> (evals)
two
$>$ (eval (list 'cdr (car '((quote (a.b)) c))))
b


## apply

- apply takes a function \& a list of arguments for it \& returns the result of applying the function to them $>\left(\right.$ apply $+{ }^{\prime}\left(\begin{array}{lll}1 & 2 & 3\end{array}\right)$ )
6
- It can be given any number of arguments, so long as the last is a list:
$>$ (apply + 12 ' $\begin{aligned} & 3 \\ & 4\end{aligned}$ 5) )
15
- A simple version of apply could be written as (define (apply flist) (eval (cons flist)))


## lambda

- The define special form creates a function and gives it a name
- However, functions don't have to have names, and we don't need to use define to create them
- The primitive way to create functions is to use the lambda special form
- These are often called lambda expressions, e.g. (lambda (x) (+ x 1))


## lambda expression

- A lambda expression is a list of the symbol lambda, followed by a list of parameters, followed by a body of one or more expressions:
$>$ (define $f($ lambda ( x ) $(+\mathrm{x} 2)$ ))
$>f$
\#[proceedure:f](proceedure:f)
$>$ (f 100)
102
$>((\operatorname{lambda}(\mathrm{x})(+\mathrm{x} 2)) 100)$ 102

| define vs. define |  |
| :---: | :---: |
| $\begin{aligned} & \text { (define (add2 } \mathrm{x} \text { ) } \\ & (+\mathrm{x} 2)) \end{aligned}$ | - The define special form comes in two varieties |
| $\begin{aligned} & \text { (define add2 } \\ & \quad(\text { lambda }(x)(+x 2))) \end{aligned}$ | The three expressions to the right are entirely equivalent |
| (define add2 \#f) (set! add2 (lambda (x) (+ x 2))) | . The first define form is just more familiar and convenient when defining a function |

## lambdas in other languages

- Lambda expressions are found in many modern languages, e.g., Python:
>>> $f=$ lambda $x, y: x^{*} x+y$
>>> f
<function <lambda> at 0x10048a230>
>> $\mathrm{f}(2,3)$
7
>> (lambda $\left.x, y: x^{*} x+y\right)(2,3)$
7


## Lambda expression

- lambda is a special form
- When evaluated, it creates a function and returns a reference to it
- The function does not have a name
- A lambda expression can be the first element of a function call:
> ( (lambda
( x$)(+\mathrm{x}$ 100)) 1) 101
- Other languages like python and javascript have adopted the idea


## More map examples

$>$ (map cons '(abc)'(123))
((a.1) (b. 2) (c. 3))
$>($ map (lambda (x) (+ x 10)) '(1 223$)$ )
(11 12 13)
$>($ map + '(1 23 3) '(45))
map: all lists must have same size; arguments were:
\#[procedure:+](procedure:+) (1 2 3) (4 5)
$===$ context $===$
/Applications/PLT/collects/scheme/private/misc.ss:7 4:7

## Defining map

Defining a simple "one argument" version of map is easy
(define (map1 func list)

```
(if (null? list)
            null
            (cons (func (first list))
                    (map1 func (rest list)))))
```


## Define Lisp's every and some

- every and some take a predicate and one or more sequences
- When given just one sequence, they test whether the elements satisfy the predicate
> (every odd? '(1 3 5)) \#t
$>$ (some even? '(1 223 3))
\#t
- If given $>1$ sequences, the predicate takes as many args as there are sequences and args are drawn one at a time from them:

$$
>\left(\text { every > '(1 } 33 \text { 5) '(0 } 2 \begin{array}{ll}
0 & 4
\end{array}\right) \text { ) }
$$

\#t

## Defining every is easy

(define (every1 f list)
;; note the use of the and function (if (null? list)
\#t
(and (f (first list))
(every1 f(rest list)))))

## Define some similarly

```
(define (some1 f list)
(if (null? list)
    #f
    (or (f (first list))
        (some1 f (rest list)))))
```


## Will this work?

- You can prove that $P$ is true for some list element by showing that it isn't false for every one
- Will this work?
$>$ (define (some1 flist)
(not (every1 (lambda (x) (not (fx))) list)))
> (some1 odd? '(2 467 8))
\#t
> (some1 (lambda (x) (> x 10)) '(4 810 12))
\#t


## filter

(filter <f> <list>) returns a list of the elements of <list> which satisfy the predicate <f>
> (filter odd? ‘(0 1234 5))
(135)
> (filter (lambda (x) (> x 98.6))
‘(101.1 98.698 .199 .4 102.2))
(101.1 99.4 102.2)

## Example: filter

(define (filter1 func list)
;; returns a list of elements of list where funct is true (cond ((null? list) null) ((func (first list)) (cons (first list) (filter1 func (rest list)))) (\#t (filter1 func (rest list)))))
> (filter1 even? '(1 23456 7))
(2 4 6)

## Example: filter

- Define integers as a function that returns a list of integers between a min and max
(define (integers min max)
(if (> min max)
null
(cons min (integers (add1 min) max))))
- Do prime? as a predicate that is true of prime numbers and false otherwise
$>$ (filter prime? (integers 2 20) )
(2 357111317 19)


## Here's another pattern

- We often want to do something like sum the elements of a sequence
(define (sum-list I)
(if (null? I)
0
(+ (first I) (sum-list (rest I)))))
- Other times we want their product
(define (multiply-list I)
(if (null? I)
1
(* (first I) (multiply-list (rest I)))))


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## Example: reduce

- Reduce takes (i) a function, (ii) a final value and (iii) a list of arguments

Reduce of,+ 0 , (v1 v2 v3 ... vn) is just $\mathrm{V} 1+\mathrm{V} 2+\mathrm{V} 3+\ldots \mathrm{Vn}+0$

- In Scheme/Lisp notation:
$>$ (reduce + 0 ‘(1 234 5))
15
(reduce * 1 '(1 234 5))
120


## Example: reduce

(define (reduce function final list)
(if (null? list)
final
(function
(first list)
(reduce function final (rest list)))))

```
(define (sum-list list)
Using reduce
    ;; returns the sum of the list elements
    (reduce + 0 list))
(define (mul-list list)
    ;; returns the sum of the list elements
    (reduce * }1\mathrm{ list))
(define (copy-list list)
    ;; copies the top level of a list
    (reduce cons '() list))
(define (append-list list)
    ;; appends all of the sublists in a list
    (reduce append '() list))
```


## The roots of mapReduce

- MapReduce is a software framework developed by Google for parallel computation on large datasets on computer clusters
- It's become an important way to
 exploit parallel computing using conventional programming languages and techniques.
- See Apache's Hadoop for an open source version

- The framework was inspired by functional programming's map, reduce \& side-effect free programs


## Defining compose

- Here's compose for two functions in Scheme (define (compose2 fg) (lambda (x) (f (g x))))
- Note that compose calls lambda which returns a new function that applies $f$ to the result of applying $g$ to $x$
- We'll look at how the variable environments work to support this in the next topic, closures
- But first, let's see how to define a general version of compose taking any number of args


## Function composition

- Math notation: $g$ oh is a composition of functions $g$ and $h$
- If $f=g \bullet h$ then $f(x)=g(h(x)$
- Composing functions is easy in Scheme
> compose
\#[procedure:compose](procedure:compose)
$>\left(\right.$ define $(\mathrm{sq} \mathrm{x})\left({ }^{*} \mathrm{x} x\right)$ )
$>$ (define $($ dub $x)(* \times 2)$ )
$>$ (sq (dub 10))
400
$>($ dub $(\mathrm{sq} 10))$
200
$>$ (define sd (compose sq dub))
$>$ (sd 10)
400
> ((compose dub sq) 10 )
200


## Functions with any number of args

- Defining functions that takes any number of arguments is easy in Scheme
(define (foo . args) (printf "My args: ~a\n" args)))
- If the parameter list ends in a symbol as opposed to null (cf. dotted pair), then it's value is the list of the remaining arguments' values (define (f x y . more-args) ...)
(define (map f. lists) ... )


## Compose in Scheme

(define (compose . FS)
;; Returns the identity function if no args given
(if (null? FS)
(lambda (x) x)
(lambda (x) ((first FS) ((apply compose (rest FS)) x)))))
; examples
(define (add-a-bang str) (string-append str "!"))
(define givebang
(compose string->symbol add-a-bang symbol->string))
(givebang 'set) ; ===> set!
; anonymous composition
((compose sqrt negate square) 5) ; ===> 0+5i

## A general every

- We can easily re-define other functions to take more than one argument
(define (every fn . args)
(cond ((null? args) \#f)
((null? (first args)) \#t)
((apply fn (map first args))
(apply every fn (map rest args)))
(else \#f)))
- (every > '(1 2 3 3) '(0 2 3)) $=>$ \#t
- (every > '(1 2 3) '(0 20 3)) $=>$ \#f


## Functional Programming Summary

- Lisp is the archetypal functional programming language
- It treated functions as first-class objects and uses the same representation for data \& code
- The FP paradigm is a good match for many problems, esp. ones involving reasoning about or optimizing code or parallel execution
- While no pure FP languages are considered mainstream, many PLs support a FP style

