## Artificial Intelligence class 3: Search (Ch. 3.1-3.3)



## Bookkeeping

- Final reminder: readings, HW, etc. are on the class schedule tiny.cc/ai-schedule
- Final reminder: the class Discord is available and useful tiny.cc/ai-discord
- HW1 is out, please verify that you can find it


## Bits From Last Time

- Sequential: Require memory of past actions to determine next best action
- Or: current action can influence all future actions
- Single- vs. multi-agent: Is "your" agent the only one affecting the world?

- Episodic: A series of one-shot actions
- Only the current percept(s) are relevant
- Sensing/acting in episode(t) is independent of episode(t-1)


## Some Examples

| Agent Type | Performance <br> Measure | Environment | Actuators | Sensors |
| :---: | :---: | :---: | :---: | :---: |
| Robot soccer <br> player | Winning game, <br> goals <br> for/against | Field, ball, <br> own team, <br> other team, <br> own body | Devices (e.g., <br> legs) for <br> locomotion <br> and kicking | Camera, touch <br> sensors, <br> accelerometers, <br> orientation <br> sensors, <br> wheel/joint <br> encoders |
| Internet <br> book-shopping <br> agent | Obtain <br> requested/ <br> Interesting <br> books, <br> minimize <br> expenditure | Internet | Follow link, <br> enter/submit <br> data in fields, <br> display to user | Web pages, <br> user requests |


| Task <br> Environment | Observable | Deterministic | Episodic | Static | Discrete | Agents |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Robot <br> soccer | Partially | Stochastic | Sequential | Dynamic | Continuous | Multi |
| Internet <br> book- <br> shopping | Partially | Deterministic | Sequential | Static | Discrete | Single |

## Environment

## Pre-Reading: Questions?

- Search (a.k.a. state-space search)
- Concepts:
- Initial state
- Transition model
- State space graph
- Step cost
- Goal test (cf. goal)
- Actions
- Path cost
- Solution / optimal solution
- Open-loop/closed-loop systems
- Expanding vs. generating a state
- The frontier (a.k.a. open list)

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## What's a "State"?

- The current value of everything in the agent's "world"
- "State space": all possible states
- Everything in the problem representation
- Values of all parameters at a particular point in time
- Examples:
- Chess board: $8 \times 8$ grid, location of all pieces
- Tic-tac-toe: $3 \times 3$ grid, whether each is $\mathrm{X}, \mathrm{O}$, or open
- Robot soccer: Location of all players, location of ball, possibly last known trajectory of all players (if sequential)
- Travel: Cities, distances between cities, agent's current city


## Today's Class

- Representing states and operators
- Example problems
- Generic state-space search algorithm
- Everything in Al comes down to search.
- Goal: understand search, and understand why.

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## Why Search?

- Traditional (non-AI) problems are likely tractable.
- Either they can be solved by listing all possible states...



...
- Tic-tac-toe: $3^{9}=19,683$ states ( 3 values for each cell, nine cells)*
- Small enough that a computer can explore all possible choices during play
- Or there's a mechanical approach to finding a solution

$$
345,781,000 \times 234,567,431,000
$$

- Can't memorize the space of answers, but you don't need to
* Of course, there are fewer valid states


## Why Search? (2)

- "Intelligent" problems are usually intractable.
- Either the state space is too large to enumerate...

- We don't know what a good solution is until we find it...
- Or, somehow, we have more
 states than we can explore.

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## Why Search? (3)

- We can't search intractable problems exhaustively, so we must consider them cleverly.
- Understanding the problem space is the first step.



## Search:The Core Idea

- For any problem:
- World is (always) in some state
- Agents take actions, which change the state
- We need a sequence of actions that gets the world into a particular goal state.
- To find it, we search the space of actions and states.

- Searching is not (always) the same as doing!


## Building Goal-Based Agents

- To build a goal-based agent we need to decide:
- What is the goal to be achieved?
- What are the possible actions?
- What relevant information must be encoded...
- To describe the state of the world?
- To describe the available transitions?
- To solve the problem?



## What is the Goal?

- A situation we want to achieve
- A set of properties that we want to hold
- Must define a "goal test" (a function over states)
- What does it mean to achieve it?
- Have we done so?
- Defining goals is a hard question that is rarely tackled in AI!
- Often, we assume the system designer or user will specify the goal
- For people, we stress the importance of establishing clear goals as the first step towards solving a problem.
- What are your goals?
- What problem(s) are you trying to solve?


## What Are Actions?

- Primitive actions or events:
- Make changes in the world
- In order to achieve a (sub)goal
- Actions are also known as operators or moves
- Examples:


## Low-level:

- Chess: "advance a pawn"
- Navigation: "take a step"
- Finance: "sell 10\% of stock X"


## High-level :

- Chess: "clear a path for a queen"
- Navigation: "go home"
- Finance: "sell best-return shares"


## Actions and Determinism

- In a deterministic world there is no uncertainty in an action's effects
- Current world state + chosen action fully specifies:
- Whether that action can be done in current world
- Is it applicable? (E.g.: Do I own any of stock X to sell?)
- Is it legal? (E.g.: Can't just move a pawn sideways.)
- World state after action is performed


## Representing Actions

- Actions here are:
- Discrete events
- That occur at an instant of time
- For example:
- State: "Mary is in class"
- Action "Go home"

- New state: "Mary is at home"
- There is no representation of a state where she is in between (i.e., in the state of "going home").


## Sliding Tile Puzzles

- 15-puzzles, 8-puzzles
- How do we represent states?
- How do we represent actions?
- Tile-1 moves north
- Tile-1 moves west

- Tile-1 moves east
- Tile-1 moves south
- Tile-2 moves north
- Tile-2 moves west



## Representing Actions

- Number of actions / operators depends on representation used in describing a state
- 8-puzzle:
- Could specify 4 possible moves (actions) for each of the 8 tiles:

initial state

goal state

$$
4 * 8=32 \text { operators. }
$$

- Or, could specify four moves for the "blank" square:

$$
4 \text { operators! }
$$

- Careful representation can simplify a problem!


## Representing States

- What information about the world sufficiently describes all aspects relevant to solving the goal?
- That is: what knowledge must be in a state description to adequately describe the current state of the world?
- The size of a problem is usually described in terms of the number of states that are possible
- Tic-Tac-Toe has about $3^{9}$ states.
- Checkers has about $10^{40}$ states.

- Rubik's Cube has about $10^{19}$ states.
- Chess has about $10^{120}$ states in a typical game.


## Some Example Problems

- Toy problems and micro-worlds
- 8-Puzzle
- Boat Problems
- Cryptarithmetic
- Remove 5 Sticks
- Water Jug Problem



## 8-Puzzle

- Given an initial configuration of 8 sliding numbered tiles on a $3 \times 3$ board, move the tiles in such a way so as to produce a desired goal configuration of the tiles.


Start State


Goal State

## 8-Puzzle

- State: $3 \times 3$ array describing where tiles are
- Operators: Move blank square Left, Right, Up or Down
- This is a more efficient encoding of the operators!
- Initial State: Starting configuration of the board
- Goal: Some specific board configuration



## The 8-Queens Problem

- Place eight (or N ) queens on a chessboard such that no queen can reach any other



## Boat Problems

- 1 sheep, 1 wolf, 1 cabbage, 1 boat
- Goal: Move everything across the river.

- Constraints:
- The boat can hold you plus one thing.
- Wolf can never be alone with sheep.
- Sheep can never be alone with cabbage.
- State: location of sheep, wolf, cabbage on shores and boat.
- Operators: Move ferry containing some set of occupants across the river (in either direction) to the other side.


## Remove 5 Sticks

- Given the following configuration of sticks, remove exactly 5 sticks in such a way that the remaining configuration forms exactly 3 squares.



## Some Real-World Problems

- Route finding
- Touring (traveling salesman)
- Logistics
- VLSI layout
- Robot navigation
- Learning



## Knowledge Representation Issues

- What's in a state?
- Is the color of the tiles relevant to solving an 8-puzzle?
- Is sunspot activity relevant to predicting the stock market?
- What to represent is a very hard problem!
- Usually left to the system designer to specify.
- What level of abstraction to describe the world?
- Too fine-grained and we "miss the forest for the trees"
- Too coarse-grained and we miss critical information


## Knowledge Representation Issues

- Number of states depends on:
- Representation choices
- Level of abstraction
- In the Remove-5-Sticks problem:
- If we represent individual sticks, then there are 17-choose-5 possible ways of removing 5 sticks (6188)
- If we represent the "squares" defined by 4 sticks, there are 6 squares initially and we must remove 3
- So, 6 -choose- 3 ways of removing 3 squares (20)


## Formalizing Search in a State Space

- A state space is a graph ( $\mathrm{V}, \mathrm{E}$ ):
- $V$ is a set of nodes (states)
- $E$ is a set of arcs (actions)
- Each arc is directed from a node to another node
- How does that work for 8puzzle?



## Formalizing Search in a State Space

- V : A node is a data structure that contains:
- State description
- Bookkeeping information: parent(s) of the node, name of operator that generated the node from that parent, etc.
- E: Each arc is an instance (single occurrence) of one operator.
- When operator is applied to the arc's source node (state), then
- Resulting state is associated with the arc's destination node



## Formalizing Search

- Each arc has a fixed, positive cost
- Corresponding to the cost of the operator
- What is "cost" of doing that action?
- Each node has a set of successor nodes
- Corresponding to all operators (actions) that can apply at source node's state
- Expanding a node is generating successor nodes, and adding them (and associated arcs) to the state-space graph
- We don't know all states initially - we have to apply operators and calculate the successor nodes


## Formalizing Search II

- One or more nodes are designated as start nodes
- A goal test predicate is applied to a state to determine if its associated node is a goal node



## Water Jug Problem as Search

Given a full 5-gallon jug and an empty 2-gallon
jug, the goal is to fill the 2-gallon jug with exactly one gallon of water.

State $=(x, y)$, where $x$ is the number of gallons of water in the 5 -gallon jug and $y$ is \# of gallons in the 2-gallon jug

Initial State $=(5,0)$
Goal State $=(*, 1)$
(* means any amount)

Operator table

| Name | Cond. | Transition | Effect |
| :---: | :---: | :---: | :---: |
| Empty5 | $1$ | ${ }^{(x) y) \rightarrow(0, y)}$ | Empty 5-gal. jug |
| Empty2 |  |  | Empty 2-gal. jug |
| 2to5 | $\begin{aligned} & x \leq 3 \\ & y=2 \end{aligned}$ | $2, \infty$ | D3yr 2-gal. into 5- |
| 5to2 | $\begin{aligned} & x \geq 2 \\ & y=0 \end{aligned}$ | $x-2,2)$ | gour 5-gal. into 2- |
| 5to2part | $\begin{aligned} & \mathrm{y}<2 \\ & \mathrm{x}=1 \end{aligned}$ | $(1, y) \rightarrow(0, y+1)$ | Pour partial 5-gal. into 2-gal. |




## Formalizing Search III

- A solution is a sequence of operators that is associated with a path in a state space from a start node to a goal node.
- 5to2, empty2, 5to2, empty2, 5to2part
- The cost of a solution is the sum of the arc costs on the solution path.
- If all arcs have the same (unit) cost, then the solution cost is just the length of the solution (number of steps / state transitions)


## Formalizing Search IV

- State-space search: searching through a state space for a solution by making explicit a sufficient portion of an implicit state-space graph to find a goal node
- Initially V=\{S\}, where S is the start node
- When S is expanded, its successors are generated; those nodes are added to V and the arcs are added to E
- This process continues until a goal node is found
- It isn't usually practical to represent entire space


## Formalizing Search V

- Each node implicitly or explicitly represents a partial solution path (and its cost) from start node to given node.
- In general, from a node there are many possible paths (and therefore solutions) that have this partial path as a prefix



## State-Space Search Algorithm

function general-search (problem, QUEUEING-FUNCTION)
;; problem describes start state, operators, goal test, and operator costs ;; queueing-function is a comparator function that ranks two states ;; returns either a goal node or failure
nodes $=$ MAKE-QUEUE(MAKE-NODE(problem.INITIAL-STATE))
loop
if EMPTY(nodes) then return "failure"
node $=$ REMOVE-FRONT(nodes)
if problem.GOAL-TEST(node.STATE) succeeds then return node
nodes $=$ QUEUEING-FUNCTION(nodes, EXPAND(node, problem.OPERATORS))
end

;; Note: The goal test is NOT done when nodes are generated
;; Note: This algorithm does not detect loops

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## Generation vs. Expansion

## - Selecting a state means making that node current

- Expanding the current state means applying every legal action to the current state
- Which generates a new set of nodes



## Key Procedures

- EXPAND
- Generate all successor nodes of a given node
- "What nodes can I reach from here (by taking what actions)?"
- GOAL-TEST
- Test if state satisfies goal conditions
- QUEUEING-FUNCTION
- Used to maintain a ranked list of nodes that are candidates for expansion
- "What should I explore next?"


## Algorithm Bookkeeping

- Typical node data structure includes:
- State at this node
- Parent node
- Operator applied to get to this node
- Depth of this node
- That is, number of operator applications since initial state
- Cost of the path
- Sum of each operator application so far


## Some Issues

- Search process constructs a search tree, where:
- Root is the initial state and
- Leaf nodes are nodes that are either:
- Not yet expanded (i.e., they are in the list "nodes") or
- Have no successors (i.e., they're "dead ends", because no operators can be applied, but they are not goals)
- Search tree may be infinite
- Even for small search space
- How?

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## Some Issues

- Return a path or a node depending on problem
- In 8-queens return a node
- 8-puzzle return a path
- What about Sheep \& Wolves?
- Changing definition of Queueing-Function $\rightarrow$ different search strategies
- How do you choose what to expand next?


## Evaluating Search Strategies

- Completeness:
- Guarantees finding a solution if one exists
- Time complexity:
- How long (worst or average case) does it take to find a solution?
- Usually measured in number of states visited/nodes expanded
- Space complexity:
- How much space is used by the algorithm?
- Usually measured in maximum size of the "nodes" list during search
- Optimality / Admissibility:
- If a solution is found, is it guaranteed to be optimal (the solution with minimum cost)?


## Summary

- Search is at the heart of AI.
- Formalizing states, actions, \&c. makes them searchable.


## Class Exercise

- Representing a Sudoku puzzle as a search space
- What are the states?
- What are the operators?
- What are the constraints (on operator application)?
- What is the description of the goal state?
- Let's try it!

|  | 3 |  |  |
| :--- | :--- | :--- | :--- |
|  |  |  | 1 |
| 3 |  |  |  |
|  |  | 2 |  |

## Sudoku, Naïvely

- State space: $4 \times 4$ matrix, divided into four $2 \times 2$ matrices: $A, B, C, D$, cells containing values [1-4]
- Operators:
- Put a 2 in square $<x, y>$
- Preconditions:
- <x,y> is empty
- $\langle x,(y \pm 1)>\neq 2 ;<x,(y \pm 2)>\neq 2 ; \ldots 3 x 4$
- $\langle(x \pm 1), y>\neq 2 ; \ldots<(x \pm 3), y>\neq 23$
- if $\langle x, y\rangle$ in $A$, then $3 \notin A$; ...

|  | 3 |  |  |
| :--- | :--- | :--- | :--- |
|  |  |  | 1 |
| 3 |  |  |  |
|  |  | 2 |  |

- How many operators is that? How many preconditions?
- Goal: all blocks are filled


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|  | 3 |  |  |
| :--- | :--- | :--- | :--- |
|  |  |  | 1 |
| 3 |  |  | 2 |
|  |  | 2 |  |

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|  | 3 |  |  |
| :--- | :--- | :--- | :--- |
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$X$ if $\langle x, y>$ in $A$, then $3 \notin A$; ...

|  | 3 |  |  |
| :--- | :--- | :--- | :--- |
|  |  |  | 1 |
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|  | 3 |  |  |
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- How many operators is that?
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# Artificial Intelligence <br> Uninformed Search (Ch. 3.4) 

(and a little more formalization)


## Questions?

- Bread-first, depth-first, uniform cost search
- Generation and expansion
- Goal tests
- Queueing function
- Complexity, completeness, and optimality


## Key Procedures

- EXPAND
- Generate all successor nodes of a given node
- "What nodes can I reach from here (by taking what actions)?"
- GOAL-TEST
- Test if state satisfies goal conditions

- QUEUEING-FUNCTION
- Maintain a ranked list of nodes that are expansion candidates
- "What should I explore next?"


## Uninformed vs. Informed Search

- Uninformed (aka "blind") search
- Use no information about the "direction" of the goal node(s)
- No way tell know if we're "doing well so far"
- Breadth-first, depth-first, depth-limited, uniform-cost, depth-first iterative deepening, bidirectional
- Informed (aka "heuristic") search (next class)
- Use domain information to (try to) (usually) head in the general direction of the goal node(s)
- Hill climbing, best-first, greedy search, beam search, A, A*


## Why Apply Goal Test Late?

- Why does it matter when the goal test is applied (expansion time vs. generation time)?
- Optimality and complexity of the algorithms are strongly affected!


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## Breadth-First

- Enqueue nodes in FIFO (first-in, first-out) order
- Characteristics:
- Complete (meaning?)
- Optimal (i.e., admissible) if all operators have the same cost
- Otherwise, not optimal but finds solution with shortest path length
- Exponential time and space complexity, $O\left(b^{d}\right)$, where:
- $d$ is the depth of the solution
- $b$ is the branching factor (number of children) at each node
- Takes a long time to find long-path solutions


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BFS



## BFS



## Breadth-First: Analysis

- Takes a long time to find long-path solutions
- Must look at all shorter length possibilities first
- A complete search tree of depth d where each non-leaf node has b children:
- $1+b+b^{2}+\ldots+b^{d}=\left(b^{d+1}-1\right) /(b-1)$ nodes
- Checks a lot of short-path solutions quickly


## Breadth-First: O(Example)

- $1+b+b^{2}+\ldots+b^{d}=\left(b^{d+1}-1\right) /(b-1)$ nodes
- Tree where: $\mathrm{d}=12$
- Every node at depths $0, \ldots, 11$ has 10 children ( $b=10$ )
- Every node at depth 12 has 0 children
- $1+10+100+1000+\ldots+10^{12}=\left(10^{13-1}\right) / 9=0\left(10^{12}\right)$ nodes in the complete search tree
- If BFS expands 1000 nodes/sec and each node uses 100 bytes of storage
- Will take 35 years to run in the worst case
- Will use 111 terabytes of memory

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## Depth-First (DFS)

- Enqueue nodes in LIFO (last-in, first-out) order
- That is, nodes used as a stack data structure to order nodes
- Characteristics:
- Might not terminate without a "depth bound"
- I.e., cutting off search below a fixed depth D ("depth-limited search")
- Not complete
- With or without cycle detection, and with or without a cutoff depth
- Exponential time, $O\left(b^{d}\right)$, but only linear space, $O(b d)$

Loops?
Infinite search spaces?


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DFS



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DFS



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DFS



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DFS


## Depth-First (DFS): Analysis

- DFS:
- Can find long solutions quickly if lucky
- And short solutions slowly if unlucky
- When search hits a dead end
- Can only back up one level at a time*
- Even if the "problem" occurs because of a bad operator choice near the top of the tree
- Hence, only does "chronological backtracking"
-     * Why?


## Uniform-Cost (UCS)

- Enqueue nodes by path cost:
- Let $\mathrm{g}(\mathrm{n})=$ cost of path from start node to current node $n$
- Sort nodes by increasing value of $g$
- Identical to breadth-first search if all operators have equal cost
- "Dijkstra's Algorithm" in algorithms literature
- "Branch and Bound Algorithm" in operations research literature
- Complete (*)
- Optimal/Admissible (*)
- Admissibility depends on the goal test being applied when a node is removed from the nodes list, not when its parent node is expanded and the node is first generated
- Exponential time and space complexity, $O\left(b^{d}\right)$


## Example: Path Costs



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## UCS Implementation

- For each frontier node, save the total cost of the path from the initial state to that node
- Expand the frontier node with the lowest path cost
- Equivalent to breadth-first if step costs all equal
- Equivalent to Dijkstra's algorithm in general


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## Uniform-cost search example



## Depth-First Iterative Deepening (DFID)

1. DFS to depth 0 (i.e., treat start node as having no successors)
until solution found do: DFS with depth cutoff $c_{\text {; }}$ $c=c+1$
2. Iff no solution, do DFS to depth 1 $\qquad$

- Complete
- Optimal/Admissible if all operators have the same cost
- Otherwise, not optimal, but guarantees finding solution of shortest length
- Time complexity is a little worse than BFS or DFS
- Nodes near the top of the tree are generated multiple times
- Because most nodes are near the bottom of a tree, worst case time complexity is still exponential, O(bd)


## Iterative deepening search ( $\mathrm{c}=1$ )



Nodes visited: 3

## Iterative deepening search ( $\mathrm{c}=2$ )



Nodes visited: 3+4 = 7

## Iterative deepening search ( $\mathrm{c}=3$ )



Nodes visited: $3+4+8=15$

## Iterative deepening search ( $\mathrm{c}=3$ )



Nodes visited: $3+4+8=15$

## Depth-First Iterative Deepening

- If branching factor is $b$ and solution is at depth $d$, then nodes at depth $d$ are generated once, nodes at depth $d-1$ are generated twice, etc.
- Hence $b^{d}+2 b^{(d-1)}+\ldots+d b \leq b^{d} /(1-1 / b)^{2}=O\left(b^{d}\right)$.
- If $b=4$, then worst case is 1.78 * $4^{\text {d }}$, i.e., $78 \%$ more nodes searched than exist at depth $d$ (in the worst case).
- Linear space complexity, O(bd), like DFS
- Has advantage of both BFS (completeness) and DFS (limited space, finds longer paths more quickly)
- Generally preferred for large state spaces where solution depth is unknown


## Example for Illustrating Search Strategies



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## Depth-First Search

Expanded node Nodes list

|  | $\left\{S^{0}\right\}$ |
| :--- | :--- |
| $S^{0}$ | $\left\{A^{3} B^{1} C^{8}\right\}$ |
| $A^{3}$ | $\left\{D^{6} E^{10} G^{18} B^{1} C^{8}\right\}$ |
| $D^{6}$ | $\left\{E^{10} G^{18} B^{1} C^{8}\right\}$ |
| $E^{10}$ | $\left\{G^{18} B^{1} C^{8}\right\}$ |
| $G^{18}$ | $\left\{B^{1} C^{8}\right\}$ |



Solution path found is S A G, cost 18
Number of nodes expanded (including goal node) $=5$

## Breadth-First Search

| Expanded node | Nodes list |
| :---: | :--- |
|  | $\left\{S^{0}\right\}$ |
| $S^{0}$ | $\left\{\mathrm{~A}^{3} \mathrm{~B}^{1} \mathrm{C}^{8}\right\}$ |
| $\mathrm{A}^{3}$ | $\left\{\mathrm{~B}^{1} \mathrm{C}^{8} \mathrm{D}^{6} \mathrm{E}^{10} \mathrm{G}^{18}\right\}$ |
| $\mathrm{B}^{1}$ | $\left\{\mathrm{C}^{8} \mathrm{D}^{6} \mathrm{E}^{10} \mathrm{G}^{18} \mathrm{G}^{21}\right\}$ |
| $\mathrm{C}^{8}$ | $\left\{\mathrm{D}^{6} \mathrm{E}^{10} \mathrm{G}^{18} \mathrm{G}^{21} \mathrm{G}^{13}\right\}$ |
| $\mathrm{D}^{6}$ | $\left\{\mathrm{E}^{10} \mathrm{G}^{18} \mathrm{G}^{21} \mathrm{G}^{13}\right\}$ |
| $\mathrm{E}^{10}$ | $\left\{\mathrm{G}^{18} \mathrm{G}^{21} \mathrm{G}^{13}\right\}$ |
| $\mathrm{G}^{18}$ | $\left\{\mathrm{G}^{21} \mathrm{G}^{13}\right\}$ |



Solution path found is S A G , cost 18
Number of nodes expanded (including goal node) $=7$

## Uniform-Cost Search

| Expanded node | Nodes list |
| :---: | :--- |
|  | $\left\{S^{0}\right\}$ |
| $S^{0}$ | $\left\{\mathrm{~B}^{1} \mathrm{~A}^{3} \mathrm{C}^{8}\right\}$ |
| $\mathrm{B}^{1}$ | $\left\{\mathrm{~A}^{3} \mathrm{C}^{8} \mathrm{G}^{21}\right\}$ |
| $\mathrm{A}^{3}$ | $\left\{\mathrm{D}^{6} \mathrm{C}^{8} \mathrm{E}^{10} \mathrm{G}^{18} \mathrm{G}^{21}\right\}$ |
| $\mathrm{D}^{6}$ | $\left\{\mathrm{C}^{8} \mathrm{E}^{10} \mathrm{G}^{18} \mathrm{G}^{1}\right\}$ |
| $\mathrm{C}^{8}$ | $\left\{\mathrm{E}^{10} \mathrm{G}^{13} \mathrm{G}^{18} \mathrm{G}^{21}\right\}$ |
| $\mathrm{E}^{10}$ | $\left\{\mathrm{G}^{13} \mathrm{G}^{18} \mathrm{G}^{21}\right\}$ |
| $\mathrm{G}^{13}$ | $\left\{\mathrm{G}^{18} \mathrm{G}^{21}\right\}$ |



Solution path found is S C G, cost 13
Number of nodes expanded (including goal node) $=7$

## How they Perform

- Depth-First Search:
- Expanded nodes: S A D E G
- Solution found: S A G (cost 18)
- Breadth-First Search:
- Expanded nodes: S A B C D E G
- Solution found: S A G (cost 18)

- Uniform-Cost Search:
- Expanded nodes: S A D B C E G
- Solution found: S C G (cost 13)
- This is the only uninformed search that worries about costs.
- Iterative-Deepening Search:
- nodes expanded: S SABCSADEG
- Solution found: S A G (cost 18)


## Comparing Search Strategies

|  | Complete | Optimal | Time complexity | Space complexity |
| :--- | :---: | :---: | :---: | :---: |
| Breadth first search: | yes | yes | $\mathrm{O}\left(\mathrm{b}^{\mathrm{d}}\right)$ | $\mathrm{O}\left(\mathrm{b}^{\mathrm{d}}\right)$ |
| Depth first search | no | no | $\mathrm{O}\left(\mathrm{b}^{\mathrm{m}}\right)$ | $\mathrm{O}(\mathrm{bm})$ |
| Depth limited search | if $1>=\mathrm{d}$ | no | $\mathrm{O}\left(\mathrm{b}^{1}\right)$ | $\mathrm{O}(\mathrm{bl})$ |
| depth first iterative <br> deepening search | yes | yes | $\mathrm{O}\left(\mathrm{b}^{\mathrm{d}}\right)$ | $\mathrm{O}(\mathrm{bd})$ |
| bi-directional search | yes | yes | $\mathrm{O}\left(\mathrm{b}^{\mathrm{d} / 2}\right)$ | $\mathrm{O}\left(\mathrm{b}^{\mathrm{d} / 2}\right)$ |

b is branching factor, d is depth of the shallowest solution, $m$ is the maximum depth of the search tree, 1 is the depth limit

## Avoiding Repeated States

- Ways to reduce size of state space (with increasing computational costs)
- In increasing order of effectiveness:

1. Do not return to the state you just came from.
2. Do not create paths with cycles in them.
3. Do not generate any state that was ever created before.

- Effect depends on frequency of loops in state space.
- Worst case, storing as many nodes as exhaustive search!


## State Space $\rightarrow$ An Exponentially Growing Search Space



## Bi-directional Search

- Alternate searching from
- start state $\rightarrow$ goal
- goal state $\rightarrow$ start
- Stop when the frontiers intersect.
- Works well only when there are unique start and goal states
- Requires ability to generate "predecessor" states.
- Can (sometimes) find a solution fast



## Bi-directional Search

- Alternate searching from
- start state $\rightarrow$ goal
- goal state $\rightarrow$ start
- Stop when the frontiers intersect.
- Works w $€$ unique st
- Requires "predecessor" states.
- Can (sometimes) find a solution fast

For next time: What's a real world problem where you can't generate predecessors?

## Holy Grail Search

Expanded node Nodes list
$S^{0} \quad\left\{\mathrm{C}^{8} \mathrm{~A}^{3} \mathrm{~B}^{1}\right\}$
$\left\{G^{13} A^{3} B^{1}\right\}$
$G^{13}$
$\left\{A^{3} B^{1}\right\}$


Solution path found is S C G, cost 13 (optimal)
Number of nodes expanded (including goal node) $=3$
(minimum possible!)

## Holy Grail Search

- Why not go straight to the solution, without any wasted detours off to the side?
- If we knew where the solution was we wouldn't be searching!

If only we knew where we were headed...

## 8-Puzzle Revisited

- What's a good algorithm?
- Depth-first search?
- Breadth-first search?
- Uniform-cost?
- Iterative deepening?


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## "Satisficing"

- Wikipedia: "Satisficing is ... searching until an acceptability threshold is met"
- Contrast with optimality

Another piece of
problem
definition

- Satisficable problems do not get more benefit from finding an optimal solution
- Ex: You have an A in the class. Studying for four hours will get you a 98 on the final. Studying for eight hours will get you a 100 on the final. What to do?
- A combination of satisfy and suffice
- Introduced by Herbert A. Simon in 1956

