

Chapter 5

Limitations of a Distr. System

- Lack of global clock
 - common clock ? Synchronized clocks ?
- Absence of shared memory
 - cannot obtain a “coherent” view of “global” state
 - coherence ==> state observations made at the same time.

Temporal fundamentals

- *Happened before* relation ($-->$)
 - $a --> b$ iff
 - a occurred before b in the same process
 - a is the event of sending a message in a process and b is the event of receiving the same message by another process
 - $-->$ is transitive
 - a can causally affect b if $a --> b$
 - if $\neg ((a --> b) \text{ and } (b --> a))$ then $a \parallel b$ (concurrent). a and b do not have a causal relationship.

Lamport's Logical Clocks

- Consider a “clock” C_i associated with process P_i . It is simply a process which assigns a number $C_i(a)$ to any event a in the process such that $C(a) < C(b)$ if $a \rightarrow b$
 - $C_i(a) < C_i(b)$ if a and b in the same process and $a \rightarrow b$
 - $C_i(a) < C_j(b)$ if a is `send(m)` in P_i and b is `recv(m)` in P_j
- To make the above true
 - C_i should monotonically increase between successive events within a process ($C_i = C_i + d$)
 - every message sent is stamped with the C_i of the sending process. On receipt, the receiver sets its C_j to the greater of its present value or the received timestamp ($\max(C_j, \text{timestamp} + d)$)
- This can be thought of as “virtual” time, but it moves only in response to events.

Limitations

- Since each clock can “independently” advance, we cannot in general infer happened before, and hence causality from clock value relations