What is a Database Management System?

1. Manages very large amounts of data.

2. Supports efficient access to very large amounts of data.

3. Supports concurrent access to v.l.a.d.
   ✦ Example: bank and its ATM machines.

4. Supports secure, atomic access to v.l.a.d.
   ✦ Contrast two people editing the same UNIX file — last to write “wins” — with the problem if two people deduct money from the same account via ATM machines at the same time — new balance is wrong whichever writes last.
Relational Model

- Based on tables, as:

<table>
<thead>
<tr>
<th>acct#</th>
<th>name</th>
<th>balance</th>
</tr>
</thead>
<tbody>
<tr>
<td>12345</td>
<td>Sally</td>
<td>1000.21</td>
</tr>
<tr>
<td>34567</td>
<td>Sue</td>
<td>285.48</td>
</tr>
<tr>
<td>...</td>
<td>...</td>
<td>...</td>
</tr>
</tbody>
</table>

- Today used in most DBMS’s.

- Amazing success story of theoretical ideas taken to $20B industry.

- “Big three” DBMS companies — Oracle, Informix, Sybase — are among the largest software companies in the world.

  ✦ IBM also plays, with its DB2 and other offerings.

  ✦ Microsoft challenges in this area too — cheap DBMS on the desktop?
• Relational companies also challenged by startup “object-oriented DB” companies.

✦ Conjecture: OODB is overblown; relational companies will prevail because they are moving rapidly to introduce OO concepts as needed via “object-relational” versions of their products and the SQL3 emerging standard.
Three Aspects to Studying DBMS’s

1. Modeling and design of databases.

   ♦ SQL = “intergalactic dataspeak.”

3. DBMS implementation.

CS145 = (1) + (2), while (3) is covered in CS245AB, CS347.
Why Study Design?

Sketch of structure for database needed, before implementation (say as relations in a relational DBMS).

- Allows discussion of issues in simplified terms before committing to implementation.

- Example: registrar maintains a DB of courses, students, instructors, etc. What goes into the DB? What assumptions can we make, e.g., unique instructor for a course?

Design Notations

Entity/relationship (E/R) and Object Description Language (ODL) will be covered in CS145, and are in common use.

- For each there is a semimechanical way to convert to relational designs (= sets of tables).
Entity/Relationship Model

- Diagrams similar in spirit to OO models.
- *Entity* like object, = “thing.”
- *Entity set* like class = set of “similar” entities/objects.
- *Attribute* = property of entities in an entity set, similar to fields or “instance variables.”
- In diagrams, entity set $\rightarrow$ rectangle; attribute $\rightarrow$ oval.

```
Students
  ID
  name
  phone
  height
```
Relationships

- Connect two or more entity sets.
- Represented by diamonds.

![Diagram of relationship between Students, Taking, and Courses]

Relationship Set

Think of the “value” of a relationship set as a table.

- One column for each of the connected entity sets.
- One row for each list of entities, one from each set, that are connected by the relationship.

<table>
<thead>
<tr>
<th>Students</th>
<th>Courses</th>
</tr>
</thead>
<tbody>
<tr>
<td>Sally</td>
<td>CS145</td>
</tr>
<tr>
<td>Sally</td>
<td>CS244</td>
</tr>
<tr>
<td>Joe</td>
<td>CS145</td>
</tr>
<tr>
<td>...</td>
<td>...</td>
</tr>
</tbody>
</table>
Multiway Relationships

Usually *binary* relationships (connecting two E.S.) suffice.

- However, there are some cases where three or more E.S. must be connected by one relationship.

- Example: relationship among students, courses, TA’s. Possibly, this E/R diagram is OK:

```
Students ----> Taking ----> Courses
           ^         ^
           |         |
           v         v
       Assisting
           |
           v
TAs
```
• Works in CS145, because each TA is a TA of all students. Connection student-TA is *only* via the course.

• But what if students were divided into sections, each headed by a TA?
  ✦ Then, a student in CS145 would be related to only one of the TA’s for CS145. Which one?

• Need a 3-way relationship to tell.
<table>
<thead>
<tr>
<th>Students</th>
<th>Courses</th>
<th>TAs</th>
</tr>
</thead>
<tbody>
<tr>
<td>Sally</td>
<td>CS145</td>
<td>Claire</td>
</tr>
<tr>
<td>Sue</td>
<td>CS145</td>
<td>Ankur</td>
</tr>
<tr>
<td>Joe</td>
<td>CS145</td>
<td>Claire</td>
</tr>
<tr>
<td>...</td>
<td>...</td>
<td>...</td>
</tr>
</tbody>
</table>
A Design Issue: Should Multiway Relationships be Split?

- Depends on whether it can be expressed as independent binary relationships.

- Example:

```
<table>
<thead>
<tr>
<th>Fossils</th>
<th>Species</th>
<th>Paleontologists</th>
</tr>
</thead>
<tbody>
<tr>
<td>Lucy</td>
<td>A. Afarensis</td>
<td>Johansson</td>
</tr>
<tr>
<td>...</td>
<td>...</td>
<td>...</td>
</tr>
</tbody>
</table>
```
Here, OK to split, because species related only to fossil, not to paleontologist.
A Design Issue: Attributes Vs. E.S. + Relationships

What is wrong with:

1. Makes it impossible to represent properties of species (e.g., height and weight of reference individual) or finders.

2. Difficult to represent a set of finders for a single fossil.

   ♦ Technically, E/R allows sets as the type of attributes, but it makes life difficult when we convert E/R to relational DB’s, so avoid sets as attribute types.

   • But if (1) and (2) are not a problem, e.g., we don’t want to represent attributes of species, then E.S. + attributes is simplest and best.
Can We Solve (1) by Adding Attributes?

- Introduces the great problem of **Redundancy**.
  - Now for every fossil of a species, we repeat the reference height and weight.
  - Wastes space (not so important these days).
  - Offers opportunity to scrub up, e.g., change height in only one fossil among several of the same species (will remain important for the foreseeable future).