

A Cryptanalyst's Dream

Peter Shor's Factoring Algorithm

by

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Crack RSA by finding
a superfast factoring
algorithm.

Problem. Given an integer N which is the product of two unknown primes p & q , i.e., $N = pq$, find p and q , i.e., factor N .

Shor's Algorithm (Cont.)

$$\mathbb{N} = \{0, 1, 2, \dots\} \quad \text{Natural Numbers}$$

Problem. Given a periodic function

$$f: \mathbb{N} \rightarrow \mathbb{N}$$

find period P of f

Choose a suff. large positive integer Q . Restrict f to

$$S_Q = \{0, 1, 2, \dots, Q-1\}$$

$$\text{Hence, } f: S_Q \rightarrow \mathbb{N}$$

Simplification. To avoid minor technicalities, we assume that Q is a multiple of P , i.e.,

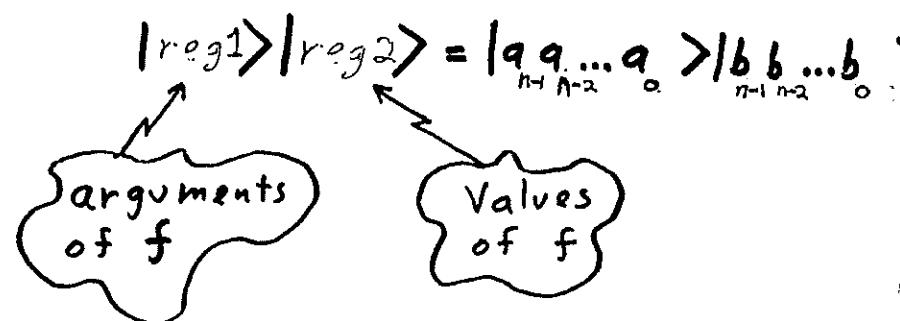
$$P | Q$$

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Shor's Algorithm (Cont.)

- Choose integer n s.t. $\begin{cases} Q < 2^n \\ \text{Max}(f) < 2^n \end{cases}$

- Construct two n -qubit registers, i.e., reg_1 and reg_2 .



Convention:

$$|a_{n-1}a_{n-2}\dots a_0\rangle = \left| \sum_{j=0}^{n-1} a_j 2^j \right\rangle$$

For example,

$$|1011\rangle = |23\rangle$$

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Shor's Algorithm (Cont.)

$$|\text{Reg1}\rangle |\text{Reg2}\rangle = \sum_{x=0}^{Q-1} |x\rangle |f(x)\rangle$$

$$= \sum_{x_0=0}^{P-1} \sum_{x_1=0}^{\frac{Q}{P}-1} |Px_1 + x_0\rangle |f(Px_1 + x_0)\rangle$$

$$= \sum_{x_0=0}^{P-1} \sum_{x_1=0}^{\frac{Q}{P}-1} |Px_1 + x_0\rangle |f(x_0)\rangle$$

$$= \sum_{x_0=0}^{P-1} \left(\sum_{x_1=0}^{\frac{Q}{P}-1} |Px_1 + x_0\rangle \right) |f(x_0)\rangle$$

Shor's Algorithm (Cont.)

$$|\text{Reg1}\rangle |\text{Reg2}\rangle = \sum_{x_0=0}^{P-1} \left(\sum_{x_1=0}^{\frac{Q}{P}-1} |Px_1 + x_0\rangle \right) |f(x_0)\rangle$$

Step 3. Measure Reg2.

$$j_0 \in \{0, 1, 2, \dots, P-1\}$$

$\text{prob}(x_0 = j_0) = \frac{1}{P}$

Whoosh!

$$|\text{Reg1}\rangle |\text{Reg2}\rangle = \sum_{x_1=0}^{\frac{Q}{P}-1} |Px_1 + j_0\rangle |f(j_0)\rangle$$

$$|\text{Reg1}\rangle |\text{Reg2}\rangle = \sum_{\lambda=0}^{P-1} \omega^{j_0 \lambda \frac{Q}{P}} |\lambda \frac{Q}{P}\rangle |f(j_0)\rangle$$

Step 5. Measure Reg1.

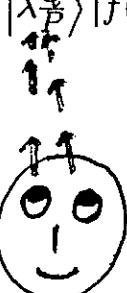
$$\lambda_0 \in \{0, 1, 2, \dots, P-1\}$$

Prob($\lambda = \lambda_0$) = $\frac{1}{P}$
Whoosh!

$$|\text{Reg1}\rangle |\text{Reg2}\rangle = |\lambda_0 \frac{Q}{P}\rangle |f(j_0)\rangle$$

Hence, we have obtained $\lambda_0 \left(\frac{Q}{P} \right)$ for some

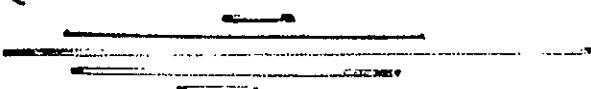
$$\lambda_0 \in \{0, 1, \dots, P-1\}$$



Repeat steps 1 thru 5 until we have obtained a large enough subset

$$\mathcal{S} \subset \{\lambda \frac{Q}{P} \mid \lambda = 0, 1, \dots, P-1\}$$

to determine $\frac{Q}{P}$, and hence P .



Note. $\frac{Q}{P}$ is obtained from \mathcal{S} by using a continued fraction expansion.

Where $\phi(P)$ denotes

Euler's totient function,

i.e., where

$$\phi(P) = \{n \in \mathbb{Z} \mid 0 < n < P \text{ and } \gcd(n, P) = 1\}$$

Reminder. The integer Q was chosen s.t.

$$N^2 \leq Q = 2^L < 2N^2$$

Observation

$\text{Prob}(y)$ is really a prob. distribution on the additive cyclic group

$$\mathbb{Z}_Q = \left(\frac{1}{Q} \mathbb{Z} \right) / \mathbb{Z} = \left\{ \frac{0}{Q}, \frac{1}{Q}, \frac{2}{Q}, \dots, \frac{Q-1}{Q} \right\} \bmod 1 ,$$

which we call the **probing quotient group**.

If a $\frac{y}{Q}$ is selected from \mathbb{Z}_Q according to Shor's prob. distribution, then the prob. is high ($\geq \frac{4}{\pi^2} \frac{\phi(P)}{P}$) that it is "close" to an element of the additive cyclic group

$$\mathbb{Z}_P = \left(\frac{1}{P} \mathbb{Z} \right) / \mathbb{Z} = \left\{ \frac{0}{P}, \frac{1}{P}, \frac{2}{P}, \dots, \frac{P-1}{P} \right\} \bmod 1 ,$$

which we call the **hidden quotient group**.

Two Norms on the Unit Circle

What Properties do these Norms Have?

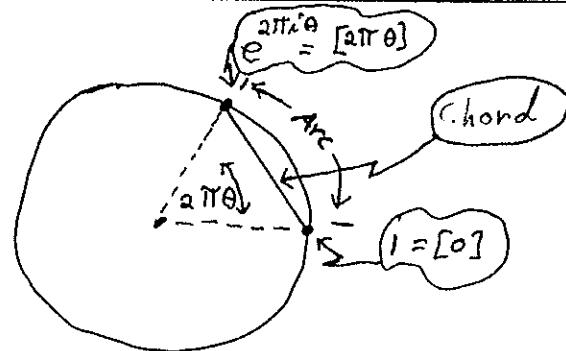
Arclength Norm

$$\text{Arc}(2\pi\theta) = 2\pi(\theta \bmod 1) = (2\pi\theta) \bmod 2\pi$$

where $\theta \bmod 1$ is the unique residue of magnitude $\leq \frac{1}{2}$, and where $(2\pi\theta) \bmod 2\pi$ is the unique residue of magnitude $\leq \pi$.

Chordal Norm

$$\text{Chord}(2\pi\theta) = 2|\sin(\pi\theta)| = |e^{2\pi i\theta} - 1|$$



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Def. A **norm** on a group G is a map

$$\|-\| : G \times G \longrightarrow \mathbf{R}$$

such that

1) $\|g\| = 0$ iff $g = \begin{cases} 1 & \text{Mult. Notation} \\ 0 & \text{Additive Notation} \end{cases}$

2) $\|g\| \geq 0$ for all $g \in G$

3) $\begin{cases} \|g_1 \cdot g_2\| \leq \|g_1\| + \|g_2\| & \text{Mult. Notation} \\ \quad \text{or} \\ \|g_1 + g_2\| \leq \|g_1\| + \|g_2\| & \text{Additive Notation} \end{cases}$

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Choose an integer

$$Q = 2^L$$

such that

$$N^2 \leq Q < 2N^2$$

Moreover, let ω be a primitive Q -th root of unity, e.g.,

$$\omega = e^{2\pi i \frac{1}{Q}}$$

Creating Shor's Prob. Distr.

Shor finds y as follows:

STEP 2.0 Initialize registers 1 and 2, i.e.,

$$|\psi_0\rangle = |\text{Reg1}\rangle |\text{Reg2}\rangle = |0\rangle |1\rangle = |00\cdots 00\rangle |00\cdots 01\rangle$$

STEP 2.1 Apply Fourier transf. F to Reg1.

$$|\psi_0\rangle = |0\rangle |1\rangle \xrightarrow{F \otimes I} |\psi_1\rangle = \frac{1}{\sqrt{Q}} \sum_{x=0}^{Q-1} |x\rangle |1\rangle$$

STEP 2.2 Apply U_f : $|x\rangle |0\rangle \mapsto |x\rangle |f(x)\rangle$

$$|\psi_1\rangle = \frac{1}{\sqrt{Q}} \sum_{x=0}^{Q-1} |x\rangle |1\rangle \xrightarrow{U_f} |\psi_2\rangle = \frac{1}{\sqrt{Q}} \sum_{x=0}^{Q-1} |x\rangle |f(x)\rangle$$

The **sole** purpose of executing STEPS 2.1 to 2.4 was to create the classical probability distribution on the set

$$\{0, 1, 2, \dots, Q - 1\}$$

given by

$$Prob(y) = \frac{\langle \Upsilon(y) | \Upsilon(y) \rangle}{Q^2},$$

where

$$|\Upsilon(y)\rangle = \sum_{x=0}^{Q-1} \omega^{xy} |f(x)\rangle = \sum_{x=0}^{Q-1} e^{\frac{2\pi i}{Q} xy} |f(x)\rangle$$

But what is $Prob(y)$?

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Recall

$$|\Upsilon(y)\rangle = \sum_{x=0}^{Q-1} \omega^{xy} |f(x)\rangle$$

Let

$$\begin{cases} Q = Pq + r, & 0 \leq r < P \\ x = Px_1 + x_0, & 0 \leq x_0 < P \end{cases}$$

Hence,

$$|\Upsilon(y)\rangle = \sum_{x_0=0}^{P-1} \omega^{x_0 y} \left(\sum_{x_1=0}^{q-1+\delta(x_0)} \omega^{Px_1 y} \right) |f(x_0)\rangle,$$

where

$$\delta(x_0) = \begin{cases} 1 & \text{if } x_0 < r \\ 0 & \text{otherwise} \end{cases}$$

$$\text{Prob}(y) = \frac{1}{Q^2} \sum_{x_0=0}^{P-1} \frac{\text{Chord}^2\left(2\pi \frac{Py}{Q}[q+\delta(x_0)]\right)}{\text{Chord}^2\left(2\pi \frac{Py}{Q}\right)}$$

$$\text{Prob}(y) \geq \frac{4}{\pi^2} \frac{1}{Q^2} \sum_{x_0=0}^{P-1} \frac{\text{Arc}^2\left(2\pi \frac{Py}{Q}[q+\delta(x_0)]\right)}{\text{Arc}^2\left(2\pi \frac{Py}{Q}\right)}$$

But ...

$$\text{Arc}(2\pi\theta) \geq \text{Chord}(2\pi\theta) \geq \frac{2}{\pi} \text{Arc}(2\pi\theta)$$

Hence,

$$\text{Prob}(y) \geq \frac{4}{\pi^2} \frac{1}{Q^2} \sum_{x_0=0}^{P-1} \frac{\text{Arc}^2\left(2\pi \frac{Py}{Q}[q+\delta(x_0)]\right)}{\text{Arc}^2\left(2\pi \frac{Py}{Q}\right)}$$

Lemma. If $\frac{y}{Q}$ is suff. "close" to a rational of the form $\frac{d}{P}$, i.e., if

$$\text{Arc}\left(2\pi \frac{y}{Q} - 2\pi \frac{d}{P}\right) \leq \frac{\pi}{Q} \left(1 - \frac{P}{Q}\right)$$

then

$$\text{Arc}\left(2\pi \frac{Py}{Q}\right) \leq \frac{\pi}{q+\delta(x_0)}$$

Also, we know that

$$\text{Arc}(2\pi\theta) \leq \frac{\pi}{|n|} \implies \text{Arc}(2\pi n\theta) = |n| \text{Arc}(2\pi\theta)$$

Recall

$$Q = Pq + r, \quad 0 \leq r < P$$

Problem

Even if $\frac{y}{Q}$ is sufficiently close to $\frac{d}{P}$, the period P can only be recovered from y if

$$\gcd(d, P) = 1$$

The probability that this happens is

$$\frac{\phi(P)}{P},$$

where $\phi(P)$ is Euler's totient function, defined by

$$\phi(P) = \#\{0 < u < P \mid \gcd(u, P) = 1\}$$

Theorem. The probability that the $\frac{y}{Q}$ produced is "sufficiently close" to a rational of the form $\frac{d}{P}$ s.t. $\gcd(d, P) = 1$ is bounded below by

$$\frac{4}{\pi^2} \frac{\phi(P)}{P} \left(1 - \frac{P}{Q}\right)^2 \gtrsim \frac{4}{\pi^2 e^\gamma} \frac{1}{\ln 2} \frac{1}{\lg \lg N} = \Omega\left(\frac{1}{\lg \lg N}\right).$$

where γ denotes Euler's constant

Second Part of Shor's Algorithm

Given a

$$0 < \frac{y}{Q} < 1$$

that is "close" to a rational of the form

$$0 < \frac{d}{P} < 1 \text{ (with } \gcd(d, P) = 1\text{)},$$

i.e., such that

$$\left| \frac{y}{Q} - \frac{d}{P} \right| \leq \frac{1}{2Q} \left(1 - \frac{P}{Q}\right),$$

how do we find the period P ?

Theorem. If

$$\left| \frac{y}{Q} - \frac{d}{P} \right| \leq \frac{1}{2P^2},$$

then $\frac{d}{P}$ is a convergent $\frac{p_k}{q_k}$ of the continued fraction expansion of $\frac{y}{Q}$.

Continued Fractions: Example

$$\begin{aligned}
 \xi = \frac{212}{97} &= 2 + \frac{18}{97} = 2 + \frac{1}{\left(\frac{97}{18}\right)} \\
 &= 2 + \frac{1}{5 + \frac{7}{18}} = 2 + \frac{1}{5 + \frac{1}{\left(\frac{18}{7}\right)}} \\
 &= 2 + \frac{1}{5 + \frac{1}{2 + \frac{4}{7}}} = 2 + \frac{1}{5 + \frac{1}{2 + \frac{1}{\left(\frac{7}{4}\right)}}} \\
 &= 2 + \frac{1}{5 + \frac{1}{2 + \frac{1}{1 + \frac{3}{4}}}} = 2 + \frac{1}{5 + \frac{1}{2 + \frac{1}{1 + \frac{1}{\left(\frac{4}{3}\right)}}}} \\
 &= 2 + \frac{1}{5 + \frac{1}{2 + \frac{1}{1 + \frac{1}{1 + \frac{1}{3}}}}} = [2, 5, 3, 1, 1, 3]
 \end{aligned}$$

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But What is a Convergent of a Continued Fraction ?

Definition. The k -th **convergent** of a continued fraction

$$[a_0, a_1, \dots, a_N]$$

is defined as

$$\frac{p_k}{q_k} = [a_0, a_1, \dots, a_k],$$

where p_k and q_k are **relatively prime**, and are given by the recursion

$$\begin{cases} p_n = a_n p_{n-1} + p_{n-2}, \\ q_n = a_n q_{n-1} + q_{n-2} . \end{cases}
 \quad \text{and} \quad
 \begin{cases} p_0 = a_0, \quad p_1 = a_1 a_0 + 1, \\ q_0 = 1, \quad q_1 = a_1, \end{cases}$$

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Why Continued Fractions?

Answer. They tell us how closely a real number can be approximated by rational numbers.

Theorem. Let ξ be a positive real number, and let d and P be integers with $P > 0$. If

$$\left| \xi - \frac{d}{P} \right| \leq \frac{1}{2P^2},$$

then

$$\frac{d}{P}$$

is a convergent of the continued fraction expansion of ξ . In other words, there exists an integer k such that

$$\frac{d}{P} = \frac{p_k}{q_k}$$